

QIC 820 Part 1 lecture 3

①

Recall: register  $X$ , associated CES  $X$ .

A quantum state is represented by a density operator.

Set of all density operators  $D(X) = \{ \rho \in \text{Pos}(X) : \text{tr} \rho = 1 \}$ .

Spectral decomposition:  $\rho = \sum_{i=1}^{\dim(X)} p(i) u_i u_i^*$ ,  $p = \text{prob vector}$   
 $\{u_i\} \subseteq X$

Def: State is pure if density operator is rank 1.

Obs:

- ①  $D(X)$  convex
- ② Extreme points of  $D(X)$ :  $\{ u u^* : u \in X, \|u\| = 1 \}$
- ③  $D(X)$  compact

To see ③,  $D(X)$  is clearly bounded, so only need to show  $D(X)$  closed, or equivalently,  $L(X) \setminus D(X)$  open.

$$L(X) \setminus D(X) = \{ A \in L(X) : A \notin \text{Herm} \} \cup \{ A \in L(X) : A \notin \text{Pos}(X) \} \\ \cup \{ A \in L(X) : \text{tr} A \neq 1 \}$$

Each of these 3 sets are open (first principle), and so is their union.



(See LN for alt proof.)

Intuitively,  $A \notin \text{Herm}$ ,  $A \notin \text{Pos}(X)$ ,  $\text{tr} A \neq 1$  are properties robust against small perturbation.

Def: Let  $X_i$ ,  $i=1, \dots, n$ , be registers,  $\rho_i \in D(X_i)$ .

Then  $\rho_1 \otimes \rho_2 \otimes \dots \otimes \rho_n \in D(X_1 \otimes X_2 \otimes \dots \otimes X_n)$ .

It is called a "product state."

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Def:  $X$  register. A measurement is specified by

①  $\Gamma =$  non-empty finite set (of outcomes)

② function  $M: \Gamma \rightarrow \text{Pos}(X)$

$$\text{s.t. } \sum_{a \in \Gamma} M(a) = \mathbb{1}_X.$$

or POVM element,  $M_a$

Each  $M(a)$  is a "measurement operator" corr to outcome  $a$ .

Axiom: If state is  $\rho \in D(X)$ , and above meas applied, then

① outcome register is in state  $\sigma = E_{a|a}$   
with prob  $p(a) = \langle M(a), \rho \rangle$

②  $X$  ceases to exist (demolition meas)

Obs: All linear functions from  $D(X)$  to prob vectors  
correspond to measurements.

Obs: We will derive non-demolition meas from demolition meas later.

Def: If  $M(a)$  is a projector for each  $a \in \Gamma$ ,  $M$  is called a projective meas

NB: Since  $\sum_{a \in \Gamma} M(a) = \mathbb{1}_X$ , the  $M(a)$ 's project onto mutually orthogonal subspaces.

Def: If  $M(a) = U_a U_a^*$  for an o.n basis  $\{U_a\}$  of  $X$

we say that the measurement is along the basis  $\{U_a\}$ .

Example = Holevo-Helstrom theorem.

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Task: Alice picks 0, 1 with prob  $p_0, p_1$ .

If outcome is  $i$ , prepares  $\rho_i$  in register  $X$ .

She gives  $X$  to Bob, who replies with  $j \in \{0, 1\}$ .

What meas maximizes  $\text{Prob}(i=j)$ ?

Lemma:  $M \in \text{Herm}(X)$ ,  $\|M\|_1 = \max \{ \text{Tr} MT : T \in \text{Herm}(X), -\mathbb{1}_X \leq T \leq \mathbb{1}_X \}$

Pf: Let  $M = \sum_k \lambda_k \chi_k \chi_k^*$  be spec decomp

$$M_{\frac{+}{-}} = \sum_{k: \lambda_k \gtrless 0} |\lambda_k| \chi_k \chi_k^*$$

$$\Pi_{\frac{+}{-}} = \sum_{k: \lambda_k > 0} \chi_k \chi_k^*$$

Then  $M = M_+ - M_-$ ,  $\|M\|_1 = \text{tr} M_+ + \text{tr} M_-$ .

(a) Let  $-\mathbb{1}_X \leq T \leq \mathbb{1}_X$ .

Define  $T_+, T_-$  similarly to  $M_{\pm}$ .

Then:  $T_+ \leq \mathbb{1}$ ,  $T_- \leq \mathbb{1}$  (omit  $X$ ).

$$\text{Tr} MT = \text{Tr} (M_+ - M_-) (T_+ - T_-)$$

$$= \underbrace{\text{Tr} M_+ T_+}_{\leq \text{Tr} M_+} + \underbrace{\text{Tr} M_- T_-}_{\leq \text{Tr} M_-} - \underbrace{\text{Tr} M_+ T_-}_{\geq 0} - \underbrace{\text{Tr} M_- T_+}_{\geq 0}$$

$$= \|M\|_1$$

(b)  $T = \Pi_+ - \Pi_-$ , all 4 ineq are equalities

$$\therefore \text{Tr} MT = \|M\|_1$$

Pf (HHT) Let Bob's meas operator be  $M_0, M_1$ . ( $M_0 = \mu_0$ )

Let  $T = M_0 - M_1$ .

$\therefore \mathbb{1} = M_0 + M_1, M_0 = \frac{1}{2} (\mathbb{1}_x + T)$

As  $0 \leq M_0 \leq \mathbb{1}, -\mathbb{1} \leq T \leq \mathbb{1}$ .

$Prob(i=j) = Prob(j=0 | i=0) \times p_0 + Prob(j=1 | i=1) \times p_1$

$= (Tr M_0 p_0) p_0 + (Tr M_1 p_1) p_1$

$= \frac{1}{2} [Tr(\mathbb{1}_x + T) p_0 p_0 + Tr(\mathbb{1}_x - T) p_1 p_1]$

$= \frac{1}{2} (1 + Tr(p_0 p_0 - p_1 p_1) T)$

$\max_M Prob(i=j) = \max_{-1 \leq T \leq 1} \frac{1}{2} (1 + Tr(p_0 p_0 - p_1 p_1) T)$

$= \frac{1}{2} (1 + \| p_0 p_0 - p_1 p_1 \|_1)$

with optimal  $T = \Pi_+ - \Pi_-$

$\uparrow$   
proj onto + espace of  $p_0 p_0 - p_1 p_1$

$M_0 = \text{proj onto } \frac{+}{-} \text{ espace of } p_0 p_0 - p_1 p_1$

## Sec 3.2 Info complete measurement: reading ex.

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### Sec 3.1.3 Product measurements

For  $n$  registers  $X_1, X_2, \dots, X_n$ , the meas

$$M: \Gamma \rightarrow \text{Pos}(X_1 \otimes \dots \otimes X_n)$$

is a product meas if  $\Gamma = \Gamma_1 \times \dots \times \Gamma_n$ , and  $\exists$  meas

$$\mu_i: \Gamma_i \rightarrow \text{Pos}(X_i)$$

$$\text{s.t. } M(a_1, \dots, a_n) = \mu_1(a_1) \otimes \mu_2(a_2) \otimes \dots \otimes \mu_n(a_n) \quad \forall a_i \in \Gamma_i, i=1, \dots, n$$

NB: when we say  $\exists$  meas  $\mu_i$ , we imply  $\sum_{j=1}^{|\Gamma_j|} \mu_i(a_j) = \mathbb{1}_{X_j}$ .

Qn: if all  $M(a_1, \dots, a_n)$  are tensor product operators, does it give a product meas?

(f. 10)

	$\pm$	
11		$\pm$
12	$\pm$	

10) 11) 12)

More on this last part of course.

### Sec 3.1.4 Channels

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Q channels transform states of one register into states of another register.

Mathematically:  $\Phi: L(X) \rightarrow L(Y)$

s.t  $\Phi$  is linear, trace-preserving, completely positive  
so that  $\mathbb{1} \otimes \Phi$  maps states to states

• trace preserving:  $\text{tr}(\Phi(A)) = \text{tr}(A)$

• completely positive:  $\forall C \in \mathbb{Z}, A \in \text{Pos}(X \otimes \mathbb{Z})$

$$\Phi \otimes \mathbb{1}_{\mathbb{Z}}(A) \in \text{Pos}(Y \otimes \mathbb{Z})$$

Physically, when  $\Phi$  is applied to  $X$  in state  $\rho$

$X$  ceases to exist, replaced by  $Y$

and state  $\rho \in D(X)$  is replaced by  $\Phi(\rho) \in D(Y)$ .

Returning to Sec 2.2,  $T(X, Y) = L(L(X), L(Y))$  (\*note linear)

$$T(X, X) =: T(X)$$

Nothing new yet:  $L(X), L(Y)$ 's are CESs.

We've learnt about linear ops in Sec 1.2.

eg Addition and scalar mult: in  $T(X, Y)$

eg.  $T(X, Y)$  is CES with dim....

eg.  $\Phi \in T(X, Y), \Phi^* \in T(Y, X) = L(L(Y), L(X))$  defined by

$$\forall A \in L(Y), B \in L(X), \langle A, \Phi(B) \rangle = \langle \Phi^*(A), B \rangle$$

eg. Tensor product (Sec 2.2.1) of  $\Phi_i: L(X_i) \rightarrow L(Y_i)$ ,  $i=1, \dots, n$  ⑦  
denoted  $\Phi_1 \otimes \Phi_2 \otimes \dots \otimes \Phi_n$   
takes  $L(X_1 \otimes X_2 \otimes \dots \otimes X_n)$  to  $L(Y_1 \otimes Y_2 \otimes \dots \otimes Y_n)$   
s.t.  $\Phi_1 \otimes \Phi_2 \otimes \dots \otimes \Phi_n (A_1 \otimes \dots \otimes A_n) = \Phi_1(A_1) \otimes \dots \otimes \Phi_n(A_n)$   
for all  $A_i \in L(X_i)$ ,  $i=1, \dots, n$ .

NB  $\Phi \in T(X, Y)$  are sometimes called superoperators  
to distinguish them from operators.

Q. What is a super-super operator? (A: maybe...)

Important superoperators and Q channels:

① Identity  $I_X: L(X) \rightarrow L(X)$   $(\mathbb{1}_{L(X)} = I_X = I)$   
 $I_X(A) = A$

linear, trace preserving, completely positive.

Also called the "noiseless channel" on  $X$ .

② Transpose  $T: L(X) \rightarrow L(X)$   
 $T(A) = A^T$

linear, trace preserving, NOT completely positive.

$$u = \sum_{i=1}^{\dim(X)} e_i \otimes e_i \quad \left( \sum_i |i\rangle\langle i| \right)$$

$$I \otimes T(u u^*) \not\geq 0$$

③ Kraus maps  $T: L(X) \rightarrow L(Y)$

$$T(A) = \sum_{k=1}^r A_k A A_k^*$$

$$\text{st. } A_k \in L(X, Y), \quad \sum_{k=1}^r A_k^* A_k = \mathbb{I}_X$$

linear, trace preserving  $\because \text{tr}(T(A)) = \sum_{k=1}^r \text{tr}(A_k A A_k^*)$

$$= \sum_{k=1}^r \text{tr}(A_k^* A_k A) = \text{tr} A$$

complete positive:  $\forall Z, \forall B \in \text{Pos}(X \otimes Z)$

$$(A_k \otimes \mathbb{I}_Z) B (A_k \otimes \mathbb{I}_Z)^* \in \text{Pos}(Y \otimes Z)$$

same when sum over  $k$ .

④ Trace:  $\text{Tr}: L(X) \rightarrow \mathbb{C}$   
 $A \mapsto \text{tr} A$

linear, trace-preserving. To see complete positivity, use an o.n basis  $\{x_i\}$

and  $\text{Tr}(A) = \sum_{i=1}^{\dim(X)} x_i^* A x_i$  st.  $\sum_{i=1}^{\dim(X)} x_i x_i^* = \mathbb{I}_X \therefore \text{Tr}$  is a Kraus map.

very important

$\text{Tr}_X$  is also a Q channel  $\forall Y$ .

$$\text{Tr}_X \otimes \mathbb{I}_Y =: \text{Tr}_X$$

Pf 1: has Kraus form  $\sum_{i=1}^{\dim(X)} (x_i^* \otimes \mathbb{I}_Y) A (x_i \otimes \mathbb{I}_Y)$

Pf 2: if  $\text{Tr}_X$  CP,  $(\text{Tr}_X \otimes \mathbb{I}_Y) \otimes \mathbb{I}_Z$  preserves positivity  $\forall Y, Z$ .

5) Measurements (Sec 6.1)

a) Non-demolition measurements / instruments

Consider a measurement on  $X$  defined by  $M: \Gamma \rightarrow \text{Pos}(X)$

$$\sum_{a \in \Gamma} \mu(a) = \mathbb{I}_X.$$

Let  $M_a \in L(X, Z)$  satisfy  $M_a^* M_a = \mu(a)$ .

(eg,  $Z = X$ ,  $M_a = \mu(a)^{\frac{1}{2}}$  function on normal ops).

$$\text{Consider } \Phi(A) = \sum_{a \in \Gamma} \underbrace{M_a A M_a^*}_{\text{in } L(Z)} \otimes \underbrace{e_a e_a^*}_{\text{in } \mathbb{C}^\Gamma} \leftarrow |a\rangle\langle a|$$

Ex: show that  $\Phi$  is linear, trace-preserving, completely-positive.

$$\text{NB } \Phi(A) = \sum_{a \in \Gamma} \underbrace{(M_a \otimes e_a)}_{\text{in } L(X, Z \otimes \mathbb{C}^\Gamma)} A (M_a \otimes e_a)^*$$

$$\begin{aligned} \text{b) } \text{Tr}_Z \Phi(A) &= \sum_{a \in \Gamma} \text{tr}(M_a A M_a^*) e_a e_a^* \\ &= \sum_{a \in \Gamma} \langle \mu_a, A \rangle e_a e_a^* = \text{meas defined by } \mu. \end{aligned}$$

Since  $\Phi$  &  $\text{Tr}_Z$  are both Q channels, so is meas defined by  $\mu$ .

(Note linearity, tr-pr, cp all preserved under composition.)

③ Partial measurements or meas one of many systems (\*) Super-important

Consider meas defined in (5.6), taking  $X$  to  $ZG$  (associated w/  $X, Z, C^{\Gamma}$ ).

Let  $Y$  be collection of all unmeasured registers.

Let  $\rho \in D(XY)$  be initial state.

Final state after measurement is:

$$\mathbb{E} \otimes I_Y (\rho) = \sum_{a \in \Gamma} (M_a \otimes I_Y) \rho (M_a^* \otimes I_Y) \otimes e_a e_a^*$$

(Sec 3.3 + Sec 6.1)

⑥ (a) Unitary channels: if  $U \in U(X)$

then  $\mathbb{E}(A) = \underbrace{U A U^{\dagger}}_{\text{Kraus map}}$  is a Q channel

⑥ (b) Mixed unitary channels if  $U_k \in U(X), k=1, \dots, r$   
Sec 6.2.3

then  $\mathbb{E}(A) = \sum_{k=1}^r p_k U_k A U_k^{\dagger}$  is a Q channel  
 $\{p_k\}$  prob vector.

⑦ Dephasing and depolarizing channels. (Sec 6.3.2)

$X$  (CS,  $\{e_a\}_{a=1}^{\dim(X)}$  fixed o.n. basis.

Dephasing channel  $\Delta(A) = \text{diag}(A)$

$$\text{i.e. } (\Delta(A))_{a,a} = A_{a,a}$$

$$(\Delta(A))_{a,b} = 0 \quad \text{if } a \neq b.$$

Depolarizing channel  $\Omega(A) = (\text{tr} A) \frac{1_X}{\dim X}$

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• If  $\mathcal{X} = (\mathbb{C}^2)^{\otimes n}$

$\sigma_0, \sigma_1, \sigma_2, \sigma_3 = \mathbb{I}_{\mathbb{C}^2}$  and Pauli  $x, y, z$  operators,  
 $(P_j)_i = \sigma_j$  on  $i$ -th qubit, tensored with  $\mathbb{I}_{\mathbb{C}^2}$  on other qubits

$$\text{then, } \Delta(A) = \frac{1}{2^n} \sum_{b_1=0}^1 \dots \sum_{b_n=0}^1 \left( \bigotimes_{i=1}^n (P_{\sigma_i})^{b_i} \right) A \left( \bigotimes_{i=1}^n (P_{\sigma_i})^{b_i} \right)^\dagger$$

↑  
Kraus maps  
∴ channels

all possible tensor product of  
 $\mathbb{I}_{\mathbb{C}^2}$  and  $\sigma_z$  on  $n$  qubits  
 as  $b_1 \dots b_n$  ranges over all  
 possible  $n$ -bit strings

$$\Downarrow$$

$$\Omega(A) = \frac{1}{4^n} \sum_{j_1=0}^3 \dots \sum_{j_n=0}^3 \left( \bigotimes_{i=1}^n (P_{\sigma_i})^{j_i} \right) A \left( \bigotimes_{i=1}^n (P_{\sigma_i})^{j_i} \right)^\dagger$$

-ranging over  $4^n$  tensor products  
 of qubit Pauli operators

• If  $\mathcal{X} = \mathbb{C}^d$ , let  $\mathbb{Z}_d = \{0, 1, \dots, d-1\}$ ,  $\omega = e^{2\pi i/d}$  (principal  $d$ -th root of unity)

$$\text{Let } X = \sum_{a \in \mathbb{Z}_d} e^{a+1} e_a^* \quad (X|a\rangle = |a+1\rangle),$$

$$Z = \sum_{a \in \mathbb{Z}_d} \omega^a e_a e_a^* \quad (Z|a\rangle = \omega^a |a\rangle).$$

Let  $W_{b,c} = X^b Z^c$ . The set  $\{W_{b,c}\}$  for  $b, c \in \mathbb{Z}_d$   
 are known as discrete Weyl operators or generalized Pauli operators.  
 or nice error basis.

Useful facts (proof as exercise):

$$- \text{Tr}(W_{b,c}) = \begin{cases} d & \text{if } b=c=0 \\ 0 & \text{otherwise.} \end{cases}$$

$$- \langle W_{g,b}, W_{c,f} \rangle = \text{Tr}(Z^{-b} X^{-g} X^c Z^f) = \text{Tr}(W_{c-g, f-b}) = \begin{cases} d & \text{if } c=g \\ & \text{and } f=b \\ 0 & \text{otherwise} \end{cases}$$

$$- ZX = WXZ$$

-  $X, Z$  generate the group  $\{W^c W_{a,b}\}_{a,b,c \in \mathbb{Z}_d}$  multiplicatively

→  $\therefore \left\{ \frac{1}{\sqrt{d}} W_{b,c} \right\}_{b,c \in \mathbb{Z}_d}$  is an o.n basis for  $L(X)$ .

Ex: check that  $\Delta(A) = \frac{1}{d} \sum_{c \in \mathbb{Z}_d} W_{0,c} A W_{0,c}^*$  (note not Hermitian but unitary).

$$\Omega(A) = \frac{1}{d^2} \sum_{b,c \in \mathbb{Z}_d} W_{b,c} A W_{b,c}^*$$

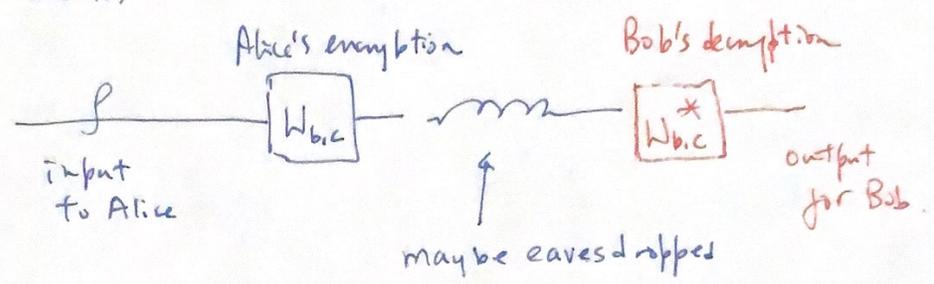
### Depolarizing channel, encryption, and teleportation

We can obtain a method to encrypt quantum states using the Kraus form for the depolarizing channel.

$$\therefore \forall A \in L(X), \quad \mathcal{L}(A) = \frac{1}{d^2} \sum_{b,c} W_{b,c} A W_{b,c}^* = \frac{\mathbb{1}_X}{d}$$

If sender Alice and receiver Bob share secret keys  $c, d$ ,

then:  $\forall \rho \in D(X)$



Without eavesdropping  $\forall b,c$ , the encryption & decryption ops cancel one another so Bob receives the input.

Without the key, an eavesdropper sees  $\frac{1}{d^2} \sum_{b,c} W_{b,c} \rho W_{b,c}^* = \frac{\mathbb{1}_X}{d}$  as the transmitted  $q$  state which is independent of the input  $\rho$ .

- This is one  $q$ . generalization of the one-time-pad to the quantum setting. It requires  $2 \log d$  key-bits of secret.

Teleportation revisited:

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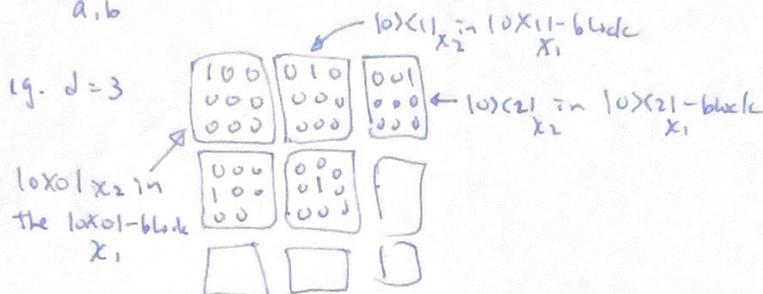
\* Lemma: for the meas  $M = \Gamma \rightarrow \text{Pos}(X)$  applied to  $X_1$ , with  $X_1 X_2$  in the maximally entangled state in  $X \otimes X$ , the post measurement state is:

$$\frac{1}{d} \sum_a |a\rangle\langle a| \otimes M(a)^T$$

Proof: assignment 1.

Def: let  $|\beta_d\rangle = \frac{1}{\sqrt{d}} \sum_{a=0}^{d-1} |a\rangle|a\rangle$  be the MES in  $\mathbb{C}^d \otimes \mathbb{C}^d$ ,

$$\beta_d = |\beta_d\rangle\langle\beta_d| = \frac{1}{d} \sum_{a,b} |a\rangle\langle b| \otimes |a\rangle\langle b|$$



Recall also the Transpose trick:  $\forall A \in L(X)$

$$A \otimes I |\beta_d\rangle = I \otimes A^T |\beta_d\rangle$$

Teleportation in d-dim:

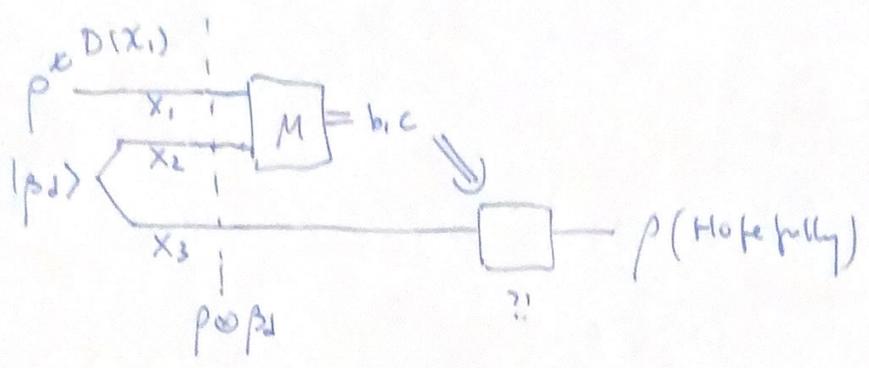
Define a measurement  $M: \mathbb{Z}_d \times \mathbb{Z}_d \rightarrow \text{Pos}(\mathbb{C}^d \otimes \mathbb{C}^d)$ ,  $\mathbb{1}_{\mathbb{C}^d} = \mathbb{1}$ .

$$M(b,c) = (W_{b,c} \otimes \mathbb{1}) \beta_d (W_{b,c}^* \otimes \mathbb{1})$$

To see that  $\sum_{b,c} M(b,c) = \mathbb{1} \otimes \mathbb{1}$ ,  $\uparrow$   
 $\text{Pos}(\mathbb{C}^d \otimes \mathbb{C}^d)$

• either note that  $\{W_{b,c} \otimes \mathbb{1} |\beta_d\rangle\}_{b,c}$  is an orthonormal basis

• or note that  $\Omega \otimes I(p) = \frac{1}{d} \mathbb{1}_{X_1} \otimes \text{tr}_{X_1} p$  for  $p \in D(X_1 \otimes X_2)$ .



• State on  $x_1 x_2 x_3$  after meas:

$$\sum_{b,c} |b,c\rangle\langle b,c| \otimes \underbrace{\text{tr}_{x_1 x_2} (M(b,c) \otimes \mathbb{1}) (\rho \otimes \beta_d)}_{//} = \text{tr}_{x_1 x_2} (\sqrt{M(b,c)} \otimes \mathbb{1}) (\rho \otimes \beta_d) (\sqrt{M(b,c)} \otimes \mathbb{1})$$

$$\text{tr}_{x_1 x_2} \left( \left[ (W_{b,c} \otimes \mathbb{1}) \beta_d (W_{b,c}^* \otimes \mathbb{1}) \right] \otimes \mathbb{1} \right) (\rho \otimes \beta_d)$$

$$\text{tr}_{x_1 x_2} \left( \beta_d (W_{b,c}^* \rho W_{b,c}) \otimes \mathbb{1} \otimes \mathbb{1} \right) (\mathbb{1} \otimes \beta_d)$$

// Lemma,  $\text{tr}_{x_1} (\beta_d (M \otimes \mathbb{1})) = M^T_{x_2}$

$$\text{tr}_{x_2} \left[ (W_{b,c}^* \rho W_{b,c})^T \otimes \mathbb{1} \right] \cdot (\beta_d)$$

//

$$(W_{b,c}^* \rho W_{b,c})^{TT} = W_{b,c}^* \rho W_{b,c}$$

∴ Bob can perform  $W_{b,c}$  if outcome  $(b,c)$  sent to him to recover  $\rho$ .

Ex: prove that  $\forall M, K$ :

$$\text{tr}_X (M \otimes \mathbb{1}) K = \text{tr}_X K (\mathbb{1} \otimes M)$$

$\uparrow \quad \uparrow \quad \uparrow$   
 $L(x) \quad L(y) \quad L(x \otimes y)$

Qn: is it true that  $\text{tr}_X K_1 K_2 = \text{tr}_X K_2 K_1$ ?