

QIC 890 / CO781 Lec 17 Nov 10, 2016

Last time:

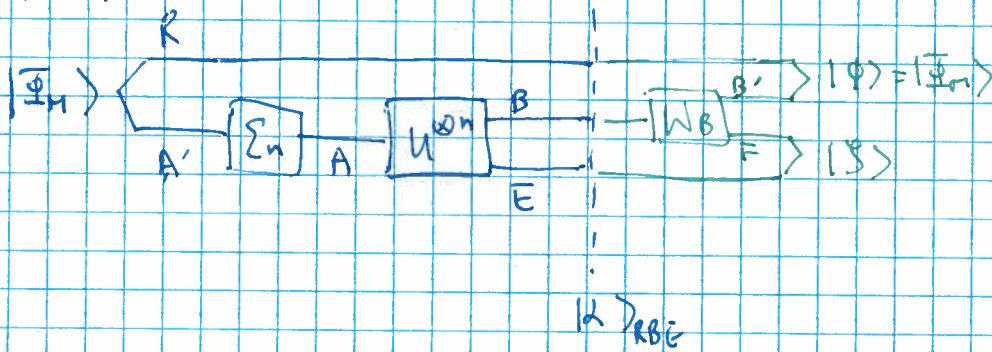
• $I_c(R>B)_p := [S(B) - S(RB)]_p = [S(B) - S(E)]_{(x)}$ where $\text{tr}_E \frac{|x\rangle\langle x|}{RBE} = \rho_{RB}$

• $Q^{(1)}(N) = \max_{|\psi\rangle} I_c(R>B)_{I \otimes N(|\psi\rangle\langle\psi|)} = \max_{|\psi\rangle} [S_B - S_E]_{(x)} = I_c(N)_{(x)}$
ISD part of N

• $Q^2(N) = \frac{1}{2} Q^{(1)}(N)$

• LSD Thm: $Q(N) = \sup_{|\psi\rangle} Q^{(1)}(N)$

• Decoupling:



$$\rho_{RE} = \left(\frac{I}{M} \right)_R \otimes \rho_E \iff \exists W_B \text{ s.t. } (W_B \otimes I_{RE}) \rho_{RBE}$$

R decoupled from E

$$\| \rho_{RB'} \otimes \rho_E \|$$

Σ_N is a code transmitting $|\Phi_M\rangle$ perfectly.

\Rightarrow
Approx version

Approx decoupling \Rightarrow a approx transmission of $|\Phi_M\rangle$

Lemma:

If $\| \rho_{RE} - \left(\frac{I}{M}\right)_R \otimes \rho_E \|_1 \leq \epsilon'$

then $\exists |\alpha\rangle_{REB}$ purifying ρ_{RE} , for $B=B'$

s.t. $\| \text{tr}_{EF} |\alpha\rangle\langle\alpha|_{REB} - |\Phi_M\rangle\langle\Phi_M|_{RB'} \|_1 \leq 2\sqrt{\epsilon'}$

Pf: (tr test \Rightarrow F \Rightarrow F of purification \Rightarrow tr test of purification \Rightarrow tr test)

Recall $F(\mu, \nu) \leq \frac{1}{2} \|\mu - \nu\|_1$ ($\forall \mu, \nu$ states)

$\therefore F(\rho_{RE}, \left(\frac{I}{M}\right)_R \otimes \rho_E) \geq 1 - \frac{\epsilon'}{2}$

since any $|\beta\rangle = u \otimes |\beta_0\rangle$, absorb u into $|\alpha\rangle$

Recall [Uhlmann] $F(\mu, \nu) = \max_{|\alpha\rangle, |\beta\rangle} |\langle\alpha|\beta\rangle| = \max_{|\alpha\rangle} |\langle\alpha|\beta_0\rangle|$
purifies μ purifies ν purifies μ fixed purification of ν

(choose $|\beta_0\rangle = |\Phi_M\rangle_{RB'} \otimes |\psi\rangle_{EF}$ where $\nu = \left(\frac{I}{M}\right)_R \otimes \rho_E$, $\mu = \rho_{RE}$)

Then: $\exists |\alpha\rangle_{REB}$ attaining the max in $F(\mu, \nu)$

$\therefore 1 - \frac{\epsilon'}{2} \leq F(\rho_{RE}, \left(\frac{I}{M}\right)_R \otimes \rho_E) = F(|\Phi_M\rangle_{RB'} \otimes |\psi\rangle_{EF}, |\alpha\rangle_{REB})$

Recall $\frac{1}{2} \|\mu - \nu\|_1 = \sqrt{1 - F(\mu, \nu)^2}$ if μ, ν pure. Apply to the 2 purifications:

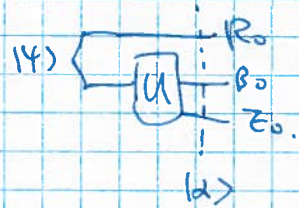
$\therefore \| |\alpha\rangle\langle\alpha|_{REB} - |\Phi_M\rangle\langle\Phi_M|_{RB'} \otimes |\psi\rangle\langle\psi|_{EF} \|_1$

$= 2 \sqrt{1 - F(|\alpha\rangle_{REB}, |\Phi_M\rangle_{RB'} |\psi\rangle_{EF})^2}$ [$(1+A) = (1+A)(1+A)$]

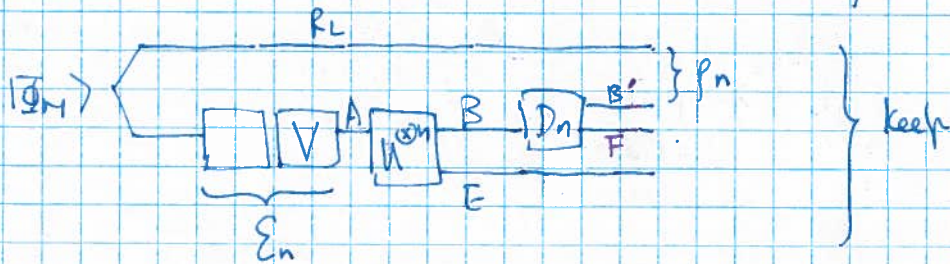
$\leq 2 \sqrt{2} \sqrt{1 - F} \leq 2\sqrt{\epsilon'}$, Finally, tracing out EF , $\| \cdot \|_1$ non-increasing.

Direct coding thm:

- Let $U =$ isometric extension of N .
- For any $|\psi\rangle_{R_0 A_0}$, let $|\omega\rangle = \underbrace{I_{R_0} \otimes U(|\psi\rangle)}_{\text{lies on } R_0 B_0 E_0}$



- We show $\exists (M, n)$ code achieving the following:



$$A = A_1 \dots A_n, \quad B = B_1 \dots B_n, \quad \bar{E} = E_1 \dots E_n$$

$$M \geq 2^{n(I(A; B))}, \quad \|\rho_n - \Phi_M\|_1 \leq \epsilon_n$$

$$r = I_c(R_0; B_0)_{|\omega\rangle}$$

- We consider a random choice of codes where

\square embeds \mathbb{C}^M into $A = A_1 \dots A_n$

\square is drawn according to the Haar measure, $V \in U(A_1 \dots A_n)$

keep on board \odot To upper bound the error ϵ_n , we show the state on

R_L & \bar{E}

(I) in ρ

$\exists D_n$ s.t.

and \bar{E} :

Requires

Then from decoupling lemma

applies

by B'

condition

called the decoupling approach to showing the direct coding thm

\odot why on average over V , R_L & \bar{E} decoupled approx

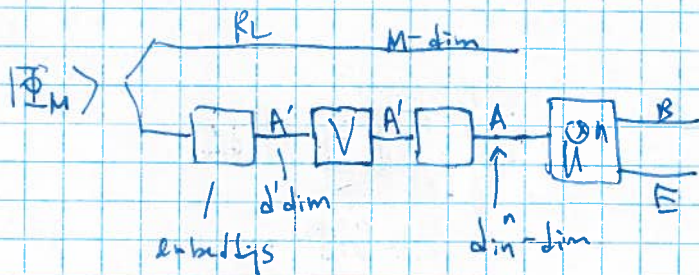
Together $\exists V$ s.t. R_L & \bar{E} decoupled enough so R_L purified by B'

resulting in the desirable ϵ_n on R_L & B' .

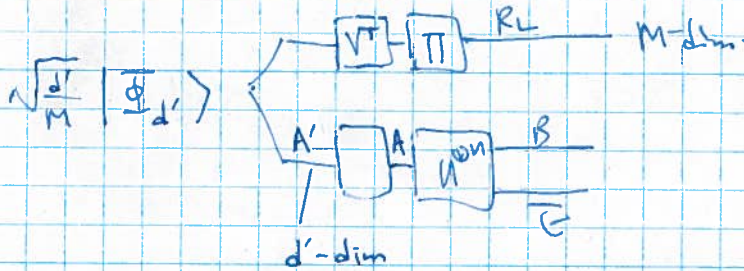
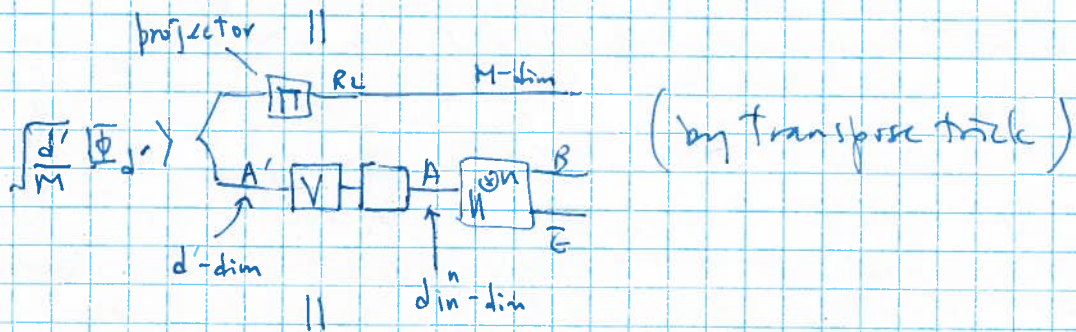
(by lemma)

- Replace N by $N^{\otimes \ell}$ and repeat the above argument gives the existence of codes attaining rate $Q^{(1)}(N)$ under the error meas. (1b) (for creating max entangled states)
- Finally, apply 9809010 / 0311037 (trimming the code space) $Q^{(1)}(N)$ is achievable for error meas (1a).

② Why on average over V , R_L & E decoupled approx:



summary
 (I)



Lemma: (0702005 Thm III)

USE Green

Let ① ρ^{SE} be max mixed on S

② $P^{S \rightarrow R}$ be a projector into subspace $R \subset S$

③ $\Psi_u^{RE} := \frac{|S|}{|R|} (P U)^S \otimes I^E \rho^{SE} (U^\dagger P)^S \otimes I^E$

Then: ① $\| \Psi_u^{RE} - \frac{|R|}{|R|} \otimes \Psi^E \|_2^2 = \text{Tr}(\Psi_u^{RE})^2 - \frac{1}{|R|} \text{Tr}(\Psi^E)^2$

② $\int du \| \Psi_u^{RE} - \frac{|R|}{|R|} \otimes \Psi^E \|_2^2 \leq \text{Tr}(\rho^{SE})^2$

③ $\int du \| \Psi_u^{RE} - \frac{|R|}{|R|} \otimes \Psi^E \|_1^2 \leq |R| \text{Tr}(\rho^{SE})^2$
|
dim R |
dim E

Pf: See supp notes on course website.

(or Watrous book)

① $\| A - B \|_2^2 = \text{Tr}(A - B)^\dagger (A - B)$
 $= \text{Tr}(A^\dagger A - B^\dagger A - A^\dagger B + B^\dagger B)$

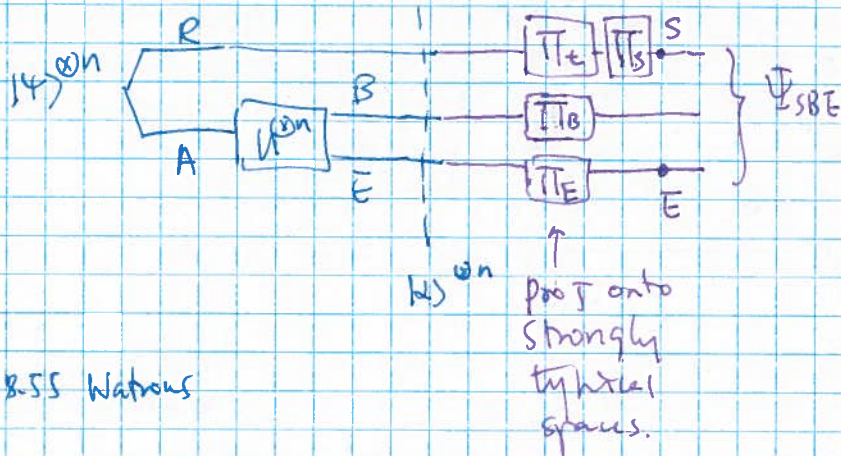
② $\int du \text{Tr}(\Psi_u^{RE})^2 = \int du \text{Tr}(\Psi_u^{RE} \otimes \Psi_u^{RE})$
|
SWAP

$\int du \text{Tr}((P \otimes P) U \otimes U [\rho^{SE} \otimes \rho^{SE}] U^\dagger \otimes U^\dagger (P \otimes P))$
|
 $= \text{Tr}(\rho^{SE} \otimes \rho^{SE}) \cdot \int du \text{Tr}(U \otimes U)$

③ $\| M \|_1^2 \leq (\text{dim } M) \| M \|_2^2$ (Cauchy-Schwarz ineq)
 $(\lambda_1 + \dots + \lambda_d)^2 \leq d (\lambda_1^2 + \dots + \lambda_d^2)$

← [This also defines the workspace]

Need a better state SBE
to apply Thm 3



proj onto a type class :

let $\Psi_R = \sum_i \lambda_i |i\rangle\langle i|$

type class t for $\Psi_R^{\otimes n}$
contains all $(z_1 \dots z_n)$
with same empirical
freq of the symbols.

$\Pi_t =$ proj onto span
of slots in the type class.
Let A_t be the space spanned.

or Lemma B.55 Watrous

Thm 4: 0702005 $\forall \epsilon > 0$

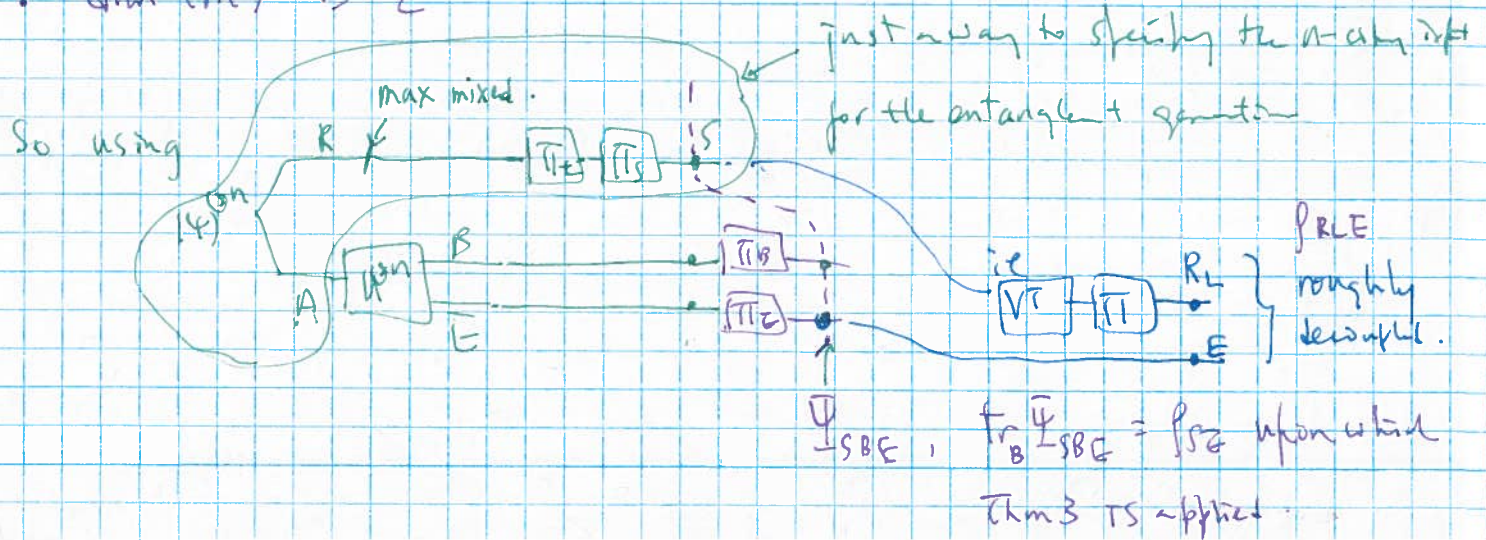
In large enough s.t. \exists type class t for $\Psi_R^{\otimes n}$ s.t.

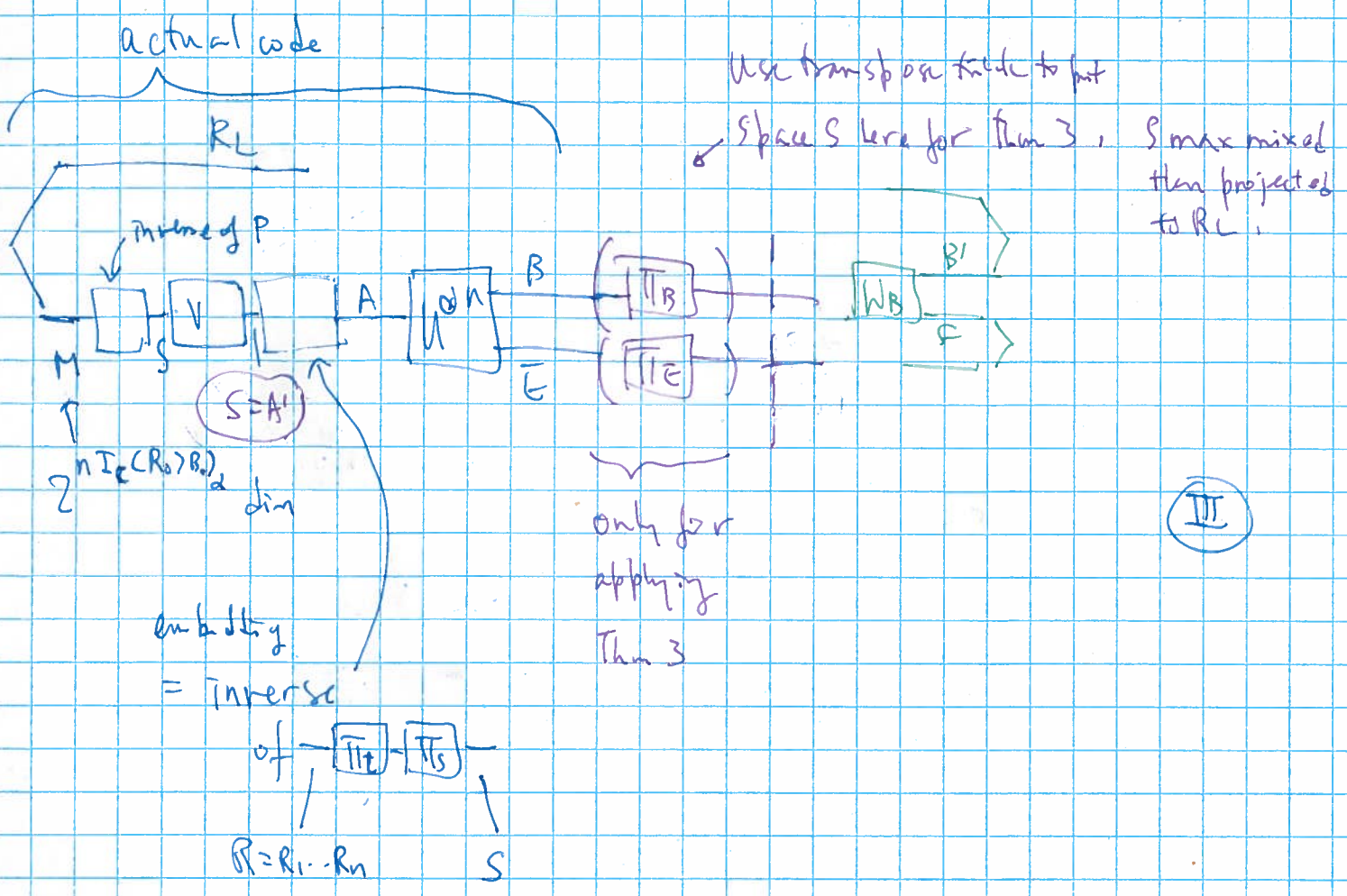
$\exists S \subset A_t, \dim(S) \geq (1 - \frac{\epsilon}{2}) \dim(A_t)$

$\left\| \frac{(\Pi_S \otimes \Pi_B \otimes \Pi_E) |\alpha\rangle^{\otimes n}}{\sqrt{\text{norm}_{\Psi}}}} - \Psi_{SBE} \right\|_1 \leq \epsilon$
w/ max mixed S

- $\text{Tr}[(\Psi_B)^2] \leq 2^{-n(S(B)_d - d)}$
- $(\dim E)_{\Psi} \leq 2^{n(S(E)_d + d)}$
- $\dim(A_t) \geq 2^{n(S(R)_d - d)}$

Ψ_{SBE} : carries the desirable
spectral / entropic properties
of $|\alpha\rangle^{\otimes n}$.





These choices allow the correct rate based on $I(S)$ simultaneously, with tracking, allows Thm 3 to apply.

$$\begin{aligned} \downarrow \text{Solv } \mathbb{T} \quad \left\| \text{PrLE} - \left(\frac{\mathbb{I}}{m} \right)_{RL} \otimes \beta_z \right\|_2^2 &\leq 2 \frac{n(\mathbb{I}_{\mathbb{C}(R_0) \times \beta_0}) - 3dn}{2} \times 2 \frac{n(\mathbb{S}(\mathbb{E}_0) \times \tau d_n)}{2} \\ &\leq 2^{-n \cdot d_n} \rightarrow 0 \end{aligned}$$

$\downarrow \exists V$ s.t. $\left\| \text{PrLE} - \left(\frac{\mathbb{I}}{m} \right)_{RL} \otimes \beta_z \right\|_2^2$ is small, say ϵ_n

\Downarrow

By lemma, can decide $\langle \mathbb{I}_m \rangle$ with trace error $2\sqrt{\epsilon_n}$.