

Optimal arrangements of $2d$ lines in \mathbb{C}^d

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THE OHIO STATE UNIVERSITY

**University of Waterloo
Virtual Algebraic Graph Theory Seminar**

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Joey Iverson
Iowa State



John Jasper
AFIT



SIM NS
FOUNDATION

tl;dr

Weak $d \times 2d$ conjecture

For every d , there exists a $d \times 2d$ **equiangular tight frame** (ETF)

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For every d , there exists a $d \times 2d$ **equiangular tight frame** (ETF)

This talk:

- ▶ two new infinite families
- ▶ conjectural construction for all d
- ▶ existence for all $d \leq 162$

Outline

1: Background

$d \times 2d$ ETF exists

1	2	3	4	5	6	7	8	9	10		12	13	14	15	16
	18	19	20	21	22	23	24	25	26	27	28		30	31	32
33	34		36	37	38		40	41	42		44	45	46		48
49	50	51	52		54	55	56	57	58		60	61	62	63	64
	66		68	69	70		72		74	75	76		78	79	80
	82		84	85	86	87	88		90	91	92		94		96
97	98	99	100		102		104		106		108		110		112
113	114	115	116	117	118		120	121	122		124		126		128
129	130		132		134	135	136		138	139	140	141	142		144
145	146	147	148		150		152		154		156	157	158	159	160

Outline

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$d \times 2d$ ETF exists

1	2	3	4	5	6	7	8	9	10	<input type="checkbox"/>	12	13	14	15	16
<input type="checkbox"/>	18	19	20	21	22	23	24	25	26	27	28	<input type="checkbox"/>	30	31	32
33	34	<input type="checkbox"/>	36	37	38	<input type="checkbox"/>	40	41	42	<input type="checkbox"/>	44	45	46	<input type="checkbox"/>	48
49	50	51	52	<input type="checkbox"/>	54	55	56	57	58	<input type="checkbox"/>	60	61	62	63	64
<input type="checkbox"/>	66	<input type="checkbox"/>	68	69	70	<input type="checkbox"/>	72	<input type="checkbox"/>	74	75	76	<input type="checkbox"/>	78	79	80
<input type="checkbox"/>	82	<input type="checkbox"/>	84	85	86	87	88	<input type="checkbox"/>	90	91	92	<input type="checkbox"/>	94	<input type="checkbox"/>	96
97	98	99	100	<input type="checkbox"/>	102	<input type="checkbox"/>	104	<input type="checkbox"/>	106	<input type="checkbox"/>	108	<input type="checkbox"/>	110	<input type="checkbox"/>	112
113	114	115	116	117	118	<input type="checkbox"/>	120	121	122	<input type="checkbox"/>	124	<input type="checkbox"/>	126	<input type="checkbox"/>	128
129	130	<input type="checkbox"/>	132	<input type="checkbox"/>	134	135	136	<input type="checkbox"/>	138	139	140	141	142	<input type="checkbox"/>	144
145	146	147	148	<input type="checkbox"/>	150	<input type="checkbox"/>	152	<input type="checkbox"/>	154	<input type="checkbox"/>	156	157	158	159	160

= previously unknown

Outline

1: Background

2: Doubling ETFs

$d \times 2d$ ETF exists

1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
<input type="checkbox"/>	18	19	20	21	22	23	24	25	26	27	28	<input type="checkbox"/>	30	31	32
33	34	35	36	37	38	39	40	41	42	43	44	45	46	47	48
49	50	51	52	<input type="checkbox"/>	54	55	56	57	58	59	60	61	62	63	64
<input type="checkbox"/>	66	67	68	69	70	71	72	<input type="checkbox"/>	74	75	76	<input type="checkbox"/>	78	79	80
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97	98	99	100	<input type="checkbox"/>	102	103	104	<input type="checkbox"/>	106	107	108	<input type="checkbox"/>	110	111	112
113	114	115	116	117	118	119	120	121	122	123	124	<input type="checkbox"/>	126	127	128
129	130	131	132	<input type="checkbox"/>	134	135	136	<input type="checkbox"/>	138	139	140	141	142	143	144
145	146	147	148	<input type="checkbox"/>	150	151	152	<input type="checkbox"/>	154	155	156	157	158	159	160

= previously unknown

Outline

1: Background

2: Doubling ETFs

3: Doubling conference graphs

$d \times 2d$ ETF exists

1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
17	18	19	20	21	22	23	24	25	26	27	28	29	30	31	32
33	34	35	36	37	38	39	40	41	42	43	44	45	46	47	48
49	50	51	52	53	54	55	56	57	58	59	60	61	62	63	64
65	66	67	68	69	70	71	72	73	74	75	76		78	79	80
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145	146	147	148	149	150	151	152		154	155	156	157	158	159	160

= previously unknown

Outline

1: Background

2: Doubling ETFs

3: Doubling conference graphs

4: 2-circulant ETFs

$d \times 2d$ ETF exists

1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
17	18	19	20	21	22	23	24	25	26	27	28	29	30	31	32
33	34	35	36	37	38	39	40	41	42	43	44	45	46	47	48
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1	2	3	4	5	6	7	8	9	10	<input type="checkbox"/>	11	12	13	14	15	16
<input type="checkbox"/> 17	18	19	20	21	22	23	24	25	26	27	28	<input type="checkbox"/> 29	30	31	32	
33	34	<input type="checkbox"/> 35	36	37	38	<input type="checkbox"/> 39	40	41	42	<input type="checkbox"/> 43	44	45	46	<input type="checkbox"/> 47	48	
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<input type="checkbox"/> 65	66	<input type="checkbox"/> 67	68	69	70	<input type="checkbox"/> 71	72	<input type="checkbox"/> 73	74	75	76	<input type="checkbox"/> 77	78	79	80	
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113	114	115	116	117	118	<input type="checkbox"/> 119	120	121	122	<input type="checkbox"/> 123	124	<input type="checkbox"/> 125	126	<input type="checkbox"/> 127	128	
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= previously unknown

Outline

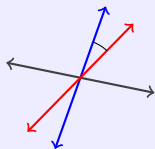
- 1: Background
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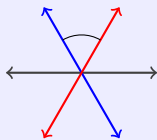
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	18	19	20	21	22	23	24	25	26	27	28		30	31	32
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Problem

Pack n lines (1-dim subspaces) in \mathbb{R}^d or \mathbb{C}^d without sharp angles



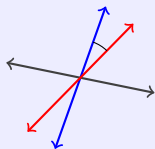
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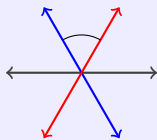
Good

Problem

Pack n lines (1-dim subspaces) in \mathbb{R}^d or \mathbb{C}^d without sharp angles



Bad



Good

- ▶ Given lines, choose unit norm reps

$$\Phi = \begin{bmatrix} | & & | \\ \varphi_1 & \cdots & \varphi_n \\ | & & | \end{bmatrix} \in \mathbb{F}^{d \times n}$$

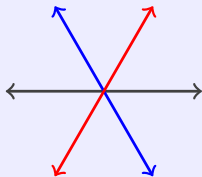


$$\cos \theta = |\langle \varphi, \psi \rangle|$$

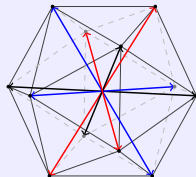
- ▶ To avoid sharp angles, minimize **coherence**

$$\mu = \max_{i \neq j} |\langle \varphi_i, \varphi_j \rangle|$$

Some optimal line packings



Real 2×3



Real 3×6

$$\frac{1}{\sqrt{2}} \begin{bmatrix} 0 & -1 & 1 & 0 & -\omega^2 & \omega & 0 & -\omega & \omega^2 \\ 1 & 0 & -1 & \omega & 0 & -\omega^2 & \omega^2 & 0 & -\omega \\ -1 & 1 & 0 & -\omega^2 & \omega & 0 & -\omega & \omega^2 & 0 \end{bmatrix}, \quad \omega = e^{2\pi i/3}$$

Complex 3×9

How do you know it's optimal?

Theorem (Welch bound)

For n unit vectors $\Phi = [\varphi_1 \ \cdots \ \varphi_n]$ in \mathbb{R}^d or \mathbb{C}^d ,

$$\mu := \max_{i \neq j} |\langle \varphi_i, \varphi_j \rangle| \geq \sqrt{\frac{n-d}{d(n-1)}}.$$

Equality holds iff Φ is an **equiangular tight frame (ETF)**:

- ▶ Equiangular: $|\langle \varphi_i, \varphi_j \rangle| = \mu$ for all $i \neq j$
- ▶ Tight frame: $\Phi\Phi^* = \text{const} \cdot I$

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equiangular tight frame (ETF) \implies optimal line packing

Certifying optimality

Example: Complex 3×9

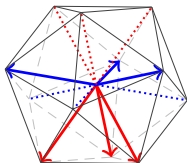
$$\Phi := \frac{1}{\sqrt{2}} \begin{bmatrix} 0 & -1 & 1 & 0 & -\omega^2 & \omega & 0 & -\omega & \omega^2 \\ 1 & 0 & -1 & \omega & 0 & -\omega^2 & \omega^2 & 0 & -\omega \\ -1 & 1 & 0 & -\omega^2 & \omega & 0 & -\omega & \omega^2 & 0 \end{bmatrix}, \quad \omega = e^{2\pi i/3}$$

$$\Phi^* \Phi = \frac{1}{2} \begin{bmatrix} 2 & -1 & -1 & -1 & -\omega & -\omega^2 & -1 & -\omega^2 & -\omega \\ -1 & 2 & -1 & -\omega^2 & -1 & -\omega & -\omega & -1 & -\omega^2 \\ -1 & -1 & 2 & -\omega & -\omega^2 & -1 & -\omega^2 & -\omega & -1 \\ -1 & -\omega & -\omega^2 & 2 & -\omega^2 & -\omega & -1 & -1 & -1 \\ -\omega^2 & -1 & -\omega & -\omega & 2 & -\omega^2 & -1 & -1 & -1 \\ -\omega & -\omega^2 & -1 & -\omega^2 & -\omega & 2 & -1 & -1 & -1 \\ -1 & -\omega^2 & -\omega & -1 & -1 & -1 & 2 & -\omega & -\omega^2 \\ -\omega & -1 & -\omega^2 & -1 & -1 & -1 & -\omega^2 & 2 & -\omega \\ -\omega^2 & -\omega & -1 & -1 & -1 & -1 & -\omega & -\omega^2 & 2 \end{bmatrix}$$

Equiangular

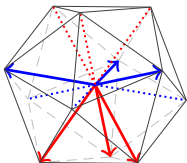
$$\Phi \Phi^* = \begin{bmatrix} 3 & 0 & 0 \\ 0 & 3 & 0 \\ 0 & 0 & 3 \end{bmatrix}$$

Tight Frame



Equiangular tight frame $\Phi = [\varphi_1 \cdots \varphi_n] \in \mathbb{F}^{d \times n}$

$$\Phi\Phi^* = \frac{n}{d}I, \quad |\langle \varphi_i, \varphi_j \rangle| = \begin{cases} 1 & i=j \\ \mu & i \neq j \end{cases}$$



Equiangular tight frame $\Phi = [\varphi_1 \cdots \varphi_n] \in \mathbb{F}^{d \times n}$

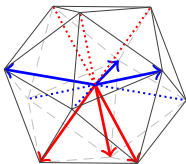
$$\Phi\Phi^* = \frac{n}{d}I, \quad |\langle \varphi_i, \varphi_j \rangle| = \begin{cases} 1 & i=j \\ \mu & i \neq j \end{cases}$$



$$G = \frac{1}{\sqrt{5}} \begin{bmatrix} \sqrt{5} & - & - & + & - & - \\ - & \sqrt{5} & - & - & + & - \\ - & - & \sqrt{5} & - & - & + \\ + & - & - & \sqrt{5} & + & + \\ - & + & - & + & \sqrt{5} & + \\ - & - & + & + & + & \sqrt{5} \end{bmatrix}$$

Gram matrix

$$G = \Phi^*\Phi \in \mathbb{F}^{n \times n}$$



Equiangular tight frame $\Phi = [\varphi_1 \cdots \varphi_n] \in \mathbb{F}^{d \times n}$

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$$\sigma(\Phi\Phi^*) = \left\{ \frac{n}{d} \right\}$$

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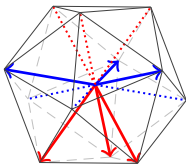


$$G = \frac{1}{\sqrt{5}} \begin{bmatrix} \sqrt{5} & - & - & + & - & - \\ - & \sqrt{5} & - & - & + & - \\ - & - & \sqrt{5} & - & - & + \\ + & - & - & \sqrt{5} & + & + \\ - & + & - & + & \sqrt{5} & + \\ - & - & + & + & + & \sqrt{5} \end{bmatrix}$$

Gram matrix

$$G = \Phi^*\Phi \in \mathbb{F}^{n \times n}$$

$$G^* = G, \quad \sigma(G) = \left\{ \frac{n}{d}, 0 \right\},$$



Equiangular tight frame $\Phi = [\varphi_1 \cdots \varphi_n] \in \mathbb{F}^{d \times n}$

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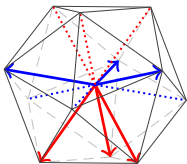
$$G = \frac{1}{\sqrt{5}} \begin{bmatrix} \sqrt{5} & - & - & + & - & - \\ - & \sqrt{5} & - & - & + & - \\ - & - & \sqrt{5} & - & - & + \\ + & - & - & \sqrt{5} & + & + \\ - & + & - & + & \sqrt{5} & + \\ - & - & + & + & + & \sqrt{5} \end{bmatrix}$$

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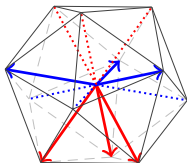
$$G^* = G, \quad \sigma(G) = \{\frac{n}{d}, 0\}, \quad |G_{ij}| = \begin{cases} 1 & i=j \\ \mu & i \neq j \end{cases}$$



Signature matrix

$$S = \frac{1}{\mu}(G - I) \in \mathbb{F}^{n \times n}$$

$$S = \begin{bmatrix} 0 & - & - & + & - & - \\ - & 0 & - & - & + & - \\ - & - & 0 & - & - & + \\ + & - & - & 0 & + & + \\ - & + & - & + & 0 & + \\ - & - & + & + & + & 0 \end{bmatrix}$$



Equiangular tight frame $\Phi = [\varphi_1 \cdots \varphi_n] \in \mathbb{F}^{d \times n}$

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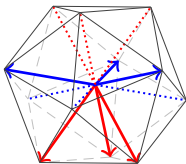
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$$S^* = S,$$

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Equiangular tight frame $\Phi = [\varphi_1 \cdots \varphi_n] \in \mathbb{F}^{d \times n}$

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Gram matrix

$$G = \Phi^*\Phi \in \mathbb{F}^{n \times n}$$

$$G^* = G, \quad \sigma(G) = \left\{ \frac{n}{d}, 0 \right\}, \quad |G_{ij}| = \begin{cases} 1 & i=j \\ \mu & i \neq j \end{cases}$$

$$G^2 = \frac{n}{d}G$$

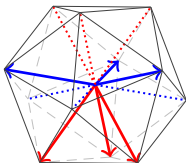


Signature matrix

$$S = \frac{1}{\mu}(G - I) \in \mathbb{F}^{n \times n}$$

$$S^* = S, \quad S^2 = \frac{1}{\mu} \left(\frac{n}{d} - 2 \right) S + (n-1)I,$$

$$S = \begin{bmatrix} 0 & - & - & + & - & - \\ - & 0 & - & - & + & - \\ - & - & 0 & - & - & + \\ + & - & - & 0 & + & + \\ - & + & - & + & 0 & + \\ - & - & + & + & + & 0 \end{bmatrix}$$



Equiangular tight frame $\Phi = [\varphi_1 \cdots \varphi_n] \in \mathbb{F}^{d \times n}$

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\Leftrightarrow

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Gram matrix

$$G = \Phi^*\Phi \in \mathbb{F}^{n \times n}$$

$$G^* = G, \quad \sigma(G) = \{\frac{n}{d}, 0\}, \quad |G_{ij}| = \begin{cases} 1 & i=j \\ \mu & i \neq j \end{cases}$$

$$G^2 = \frac{n}{d}G$$

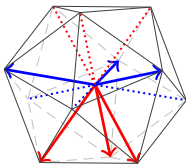
\Leftrightarrow

Signature matrix

$$S = \frac{1}{\mu}(G - I) \in \mathbb{F}^{n \times n}$$

$$S^* = S, \quad S^2 = \frac{1}{\mu}(\frac{n}{d} - 2)S + (n-1)I, \quad |S_{ij}| = \begin{cases} 0 & i=j \\ 1 & i \neq j \end{cases}$$

$$S = \begin{bmatrix} 0 & - & - & + & - & - \\ - & 0 & - & - & + & - \\ - & - & 0 & - & - & + \\ + & - & - & 0 & + & + \\ - & + & - & + & 0 & + \\ - & - & + & + & + & 0 \end{bmatrix}$$



Equiangular tight frame $\Phi = [\varphi_1 \cdots \varphi_n] \in \mathbb{F}^{d \times n}$

$$\Phi\Phi^* = \frac{n}{d}I, \quad |\langle \varphi_i, \varphi_j \rangle| = \begin{cases} 1 & i=j \\ \mu & i \neq j \end{cases}$$

$$\sigma(\Phi\Phi^*) = \{\frac{n}{d}\}$$



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$$C = \begin{bmatrix} 0 & + & + & + \\ - & 0 & + & - \\ - & - & 0 & + \\ - & + & - & 0 \end{bmatrix}$$

Conference matrix $C \in \mathbb{R}^{n \times n}$:

- ▶ $C_{ij} \in \{1, -1\}$ for every $i \neq j$
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Proof: $S := \begin{cases} iC & \text{if } C^T = -C \\ C & \text{if } C^T = C \end{cases}$

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$$\Phi = \frac{1}{\sqrt{3}} \begin{bmatrix} i\sqrt{3} & 1 & 1 & 1 \\ 0 & \sqrt{2} & \omega\sqrt{2} & \bar{\omega}\sqrt{2} \end{bmatrix}$$

$\omega = e^{2\pi i/3}$

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Theorem (Paley)

q odd prime power \implies conference matrix of order $q + 1$
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Construction: Legendre $\chi: \mathbb{F}_q \rightarrow \{0, 1, -1\}$

$$q = 3$$

a	0	1	2
$\chi(a)$	0	+	-

$$\chi(a) = \begin{cases} 0 & a = 0 \\ 1 & a \neq 0 \text{ is a square} \\ -1 & \text{else} \end{cases}$$

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$$Q = \begin{matrix} & 0 & 1 & 2 \\ \begin{matrix} 0 \\ 1 \\ 2 \end{matrix} & \begin{bmatrix} 0 & + & - \\ - & 0 & + \\ + & - & 0 \end{bmatrix} \end{matrix}$$

$$Q := [\chi(b - a)]_{a,b \in \mathbb{F}_q} \in \mathbb{R}^{q \times q}$$

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$d \times 2d$ ETF exists

1	2	3	4	5	6	7	8	9	10		12	13	14	15	16
	18	19	20	21	22	23	24	25	26	27	28		30	31	32
33	34		36	37	38		40	41	42		44	45	46		48
49	50	51	52		54	55	56	57	58		60	61	62	63	64
	66		68	69	70		72		74	75	76		78	79	80
	82		84	85	86	87	88		90	91	92		94		96
97	98	99	100		102		104		106		108		110		112
113	114	115	116	117	118		120	121	122		124		126		128
129	130		132		134	135	136		138	139	140	141	142		144
145	146	147	148		150		152		154		156	157	158	159	160

Outline

1: Background

2: Doubling ETFs

3: Doubling conference graphs

4: 2-circulant ETFs

$d \times 2d$ ETF exists

1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
<input type="checkbox"/>	18	19	20	21	22	23	24	25	26	27	28	<input type="checkbox"/>	30	31	32
33	34	35	36	37	38	39	40	41	42	43	44	45	46	47	48
49	50	51	52	<input type="checkbox"/>	54	55	56	57	58	59	60	61	62	63	64
<input type="checkbox"/>	66	67	68	69	70	71	72	<input type="checkbox"/>	74	75	76	<input type="checkbox"/>	78	79	80
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= previously unknown

Recall

Two options for conference matrix of order $n > 2$:

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Inspiration

If $C \in \mathbb{R}^{n \times n}$ is a skew-symmetric conference matrix, then so is

$$K := \begin{bmatrix} C & C+I \\ C-I & -C \end{bmatrix} \in \mathbb{R}^{2n \times 2n}$$

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Check:

$$K^T = \begin{bmatrix} C^T & C^T - I \\ C^T + I & -C^T \end{bmatrix} = \begin{bmatrix} -C & -C - I \\ -C + I & C \end{bmatrix} = -K$$

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$$\begin{aligned} KK^T &= \begin{bmatrix} C & C+I \\ C-I & -C \end{bmatrix} \begin{bmatrix} C^T & C^T - I \\ C^T + I & -C^T \end{bmatrix} \\ &= \dots \\ &= \begin{bmatrix} (2n-1)I & 0 \\ 0 & (2n-1)I \end{bmatrix} \end{aligned}$$

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Key idea

This generalizes for other ETF signature matrices

ETF Doubling Theorem (Fallon, Iverson)

$$d \times n \text{ ETF}, n \in \{2d, 2d \pm 1\} \implies n \times 2n \text{ ETF}$$

via signature matrices

$$S \mapsto \begin{bmatrix} S & S + \beta I \\ S + \bar{\beta} I & -S \end{bmatrix} =: \Sigma, \quad \beta = \beta(d, n) \in \mathbb{T}$$

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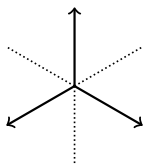
$$\begin{aligned} \Sigma^2 &= \begin{bmatrix} S & S + \beta I \\ S + \bar{\beta} I & -S \end{bmatrix} \begin{bmatrix} S & S + \beta I \\ S + \bar{\beta} I & -S \end{bmatrix} \\ &= \dots \\ &= \begin{bmatrix} (2n - 1)I & 0 \\ 0 & (2n - 1)I \end{bmatrix} \end{aligned}$$

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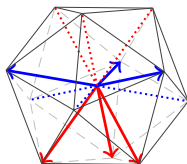
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2×3

\mapsto



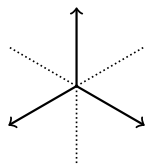
3×6

ETF Doubling Theorem (Fallon, Iverson)

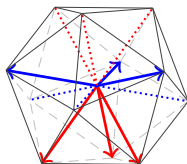
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2×3



3×6



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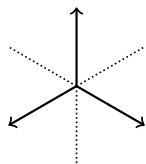
6×12

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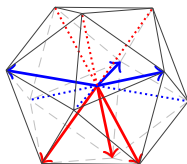
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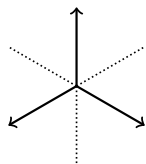
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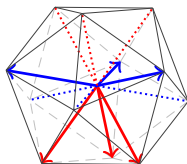
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► Sufficient: $d \times 2d$ ETF for all **odd** d

An ETF to double

Theorem (Strohmer/Renes)

Skew-conference matrix order $n \implies$ ETF size $\frac{n-2}{2} \times (n-1)$

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$$C = \begin{bmatrix} 0 & + & + & + \\ - & 0 & + & - \\ - & - & 0 & + \\ - & + & - & 0 \end{bmatrix}$$

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<input type="checkbox"/>	18	19	20	21	22	23	24	25	26	27	28	<input type="checkbox"/>	30	31	32
33	34	35	36	37	38	39	40	41	42	43	44	45	46	47	48
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= previously unknown

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$$H = \begin{bmatrix} + & + & + & + \\ - & + & + & - \\ - & - & + & + \\ - & + & - & + \end{bmatrix}$$

$$C = \begin{bmatrix} 0 & + & + & + \\ - & 0 & + & - \\ - & - & 0 & + \\ - & + & - & 0 \end{bmatrix}$$

Skew-Hadamard $H \in \mathbb{R}^{n \times n}$:

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- ▶ $HH^T = nI$
- ▶ $C^T = -C$ for $C := H - I$

$$H = \begin{bmatrix} + & + & + & + \\ - & + & + & - \\ - & - & + & + \\ - & + & - & + \end{bmatrix}$$

Skew-conference $C \in \mathbb{R}^{n \times n}$:

- ▶ $C_{ij} \in \{+, -\}$ for every $i \neq j$
- ▶ $CC^T = (n-1)I$
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$$C = \begin{bmatrix} 0 & + & + & + \\ - & 0 & + & - \\ - & - & 0 & + \\ - & + & - & 0 \end{bmatrix}$$

Skew-Hadamard $H \in \mathbb{R}^{n \times n}$:

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- ▶ $HH^T = nI$
- ▶ $C^T = -C$ for $C := H - I$

$$H = \begin{bmatrix} + & + & + & + \\ - & + & + & - \\ - & - & + & + \\ - & + & - & + \end{bmatrix}$$

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$$C = \begin{bmatrix} 0 & + & + & + \\ - & 0 & + & - \\ - & - & 0 & + \\ - & + & - & 0 \end{bmatrix}$$

skew-Hadamard $\xleftrightarrow{\pm I}$ **skew-conference:**
 $HH^T = (C + I)(C + I)^T = CC^T + (C + C^T) + I = CC^T + I$

Skew-Hadamard $H \in \mathbb{R}^{n \times n}$:

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$$\begin{aligned} \text{skew-Hadamard} & \xleftrightarrow{\pm I} \text{skew-conference:} \\ HH^T &= (C + I)(C + I)^T = CC^T + (C + C^T) + I = CC^T + I \end{aligned}$$

Skew-Hadamard conjecture

There is a skew-Hadamard matrix of every order divisible by 4

Skew-Hadamard $H \in \mathbb{R}^{n \times n}$:

- ▶ $H_{ij} \in \{+, -\}$ for every i, j
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$$H = \begin{bmatrix} + & + & + & + \\ - & + & + & - \\ - & - & + & + \\ - & + & - & + \end{bmatrix}$$

Skew-conference $C \in \mathbb{R}^{n \times n}$:

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$$C = \begin{bmatrix} 0 & + & + & + \\ - & 0 & + & - \\ - & - & 0 & + \\ - & + & - & 0 \end{bmatrix}$$

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Skew-Hadamard conjecture

There is a skew-Hadamard matrix of every order divisible by 4

Implications for $d \times 2d$ ETFs:

- ▶ all even d
- ▶ all $d \equiv 3 \pmod{4}$

Outline

- 1: Background
- 2: Doubling ETFs
- 3: Doubling conference graphs
- 4: 2-circulant ETFs

$d \times 2d$ ETF exists

1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
17	18	19	20	21	22	23	24	25	26	27	28	29	30	31	32
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145	146	147	148	149	150	151	152		154	155	156	157	158	159	160

= previously unknown

Recall

Two options for conference matrix of order $n > 2$:

- ▶ $C^T = -C, n \equiv 0 \pmod{4}$
- ▶ $C^T = C, n \equiv 2 \pmod{4}$

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Sym conference mat order $n \iff$ **Conference graph** order $n - 1$

Sym conference mat order $n \iff$ **Conference graph** order $n - 1$

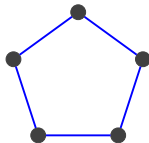
$$C = \begin{bmatrix} 0 & + & + & + & + & + \\ + & 0 & + & - & - & + \\ + & + & 0 & + & - & - \\ + & - & + & 0 & + & - \\ + & - & - & + & 0 & + \\ + & + & - & - & + & 0 \end{bmatrix}$$

Sym conference mat order $n \iff$ **Conference graph** order $n - 1$

$$C = \left[\begin{array}{c|ccccc} 0 & + & + & + & + & + \\ \hline + & 0 & + & - & - & + \\ + & + & 0 & + & - & - \\ + & - & + & 0 & + & - \\ + & - & - & + & 0 & + \\ + & + & - & - & + & 0 \end{array} \right]$$

Sym conference mat order $n \iff$ **Conference graph** order $n - 1$

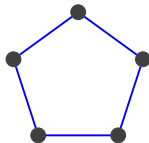
$$C = \left[\begin{array}{c|ccccc} 0 & + & + & + & + & + \\ + & 0 & + & - & - & + \\ + & + & 0 & + & - & - \\ + & - & + & 0 & + & - \\ + & - & - & + & 0 & + \\ + & + & - & - & + & 0 \end{array} \right] \mapsto A = \begin{bmatrix} 0 & 0 & 1 & 1 & 0 \\ 0 & 0 & 0 & 1 & 1 \\ 1 & 0 & 0 & 0 & 1 \\ 1 & 1 & 0 & 0 & 0 \\ 0 & 1 & 1 & 0 & 0 \end{bmatrix}, \quad B = \begin{bmatrix} 0 & 1 & 0 & 0 & 1 \\ 1 & 0 & 1 & 0 & 0 \\ 0 & 1 & 0 & 1 & 0 \\ 0 & 0 & 1 & 0 & 1 \\ 1 & 0 & 0 & 1 & 0 \end{bmatrix}$$



Sym conference mat order $n \iff$ **Conference graph** order $n - 1$

$$C = \begin{bmatrix} 0 & + & + & + & + & + \\ + & 0 & + & - & - & + \\ + & + & 0 & + & - & - \\ + & - & + & 0 & + & - \\ + & - & - & + & 0 & + \\ + & + & - & - & + & 0 \end{bmatrix} \mapsto A = \begin{bmatrix} 0 & 0 & 1 & 1 & 0 \\ 0 & 0 & 0 & 1 & 1 \\ 1 & 0 & 0 & 0 & 1 \\ 1 & 1 & 0 & 0 & 0 \\ 0 & 1 & 1 & 0 & 0 \end{bmatrix}, \quad B = \begin{bmatrix} 0 & 1 & 0 & 0 & 1 \\ 1 & 0 & 1 & 0 & 0 \\ 0 & 1 & 0 & 1 & 0 \\ 0 & 0 & 1 & 0 & 1 \\ 1 & 0 & 0 & 1 & 0 \end{bmatrix}$$

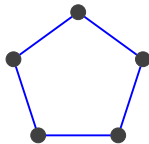
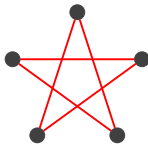
- ▶ all verts:
 $k = \frac{n-2}{2}$ neighbors
- ▶ adj verts:
 $\lambda = \frac{n-6}{2}$ common neighbors
- ▶ non-adj verts:
 $\mu = \frac{n-2}{4}$ common neighbors



Sym conference mat order $n \iff$ **Conference graph** order $n - 1$

$$C = \left[\begin{array}{c|ccccc} 0 & + & + & + & + & + \\ + & 0 & + & - & - & + \\ + & + & 0 & + & - & - \\ + & - & + & 0 & + & - \\ + & - & - & + & 0 & + \\ + & + & - & - & + & 0 \end{array} \right] \mapsto A = \begin{bmatrix} 0 & 0 & 1 & 1 & 0 \\ 0 & 0 & 0 & 1 & 1 \\ 1 & 0 & 0 & 0 & 1 \\ 1 & 1 & 0 & 0 & 0 \\ 0 & 1 & 1 & 0 & 0 \end{bmatrix}, \quad B = \begin{bmatrix} 0 & 1 & 0 & 0 & 1 \\ 1 & 0 & 1 & 0 & 0 \\ 0 & 1 & 0 & 1 & 0 \\ 0 & 0 & 1 & 0 & 1 \\ 1 & 0 & 0 & 1 & 0 \end{bmatrix}$$

- ▶ all verts:
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$$A^2 = \frac{n-6}{4}A + \frac{n-2}{4}B + \frac{n-2}{2}I$$

$$B^2 = \frac{n-2}{4}A + \frac{n-6}{4}B + \frac{n-2}{2}I$$

$$AB = BA = \frac{n-2}{4}A + \frac{n-2}{4}B$$

Conference Graph Doubling Theorem (Iverson, Jasper, M)

Sym conference mat order $n \implies$ ETF size $(n - 1) \times 2(n - 1)$

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Conference \mapsto ETF signature:

$$C = \begin{bmatrix} 0 & + & + & + & + & + \\ + & 0 & + & - & - & + \\ + & + & 0 & + & - & - \\ + & - & + & 0 & + & - \\ + & - & - & + & 0 & + \\ + & + & - & - & + & 0 \end{bmatrix}$$

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$$\alpha = \alpha(n), \beta = \beta(n) \in \mathbb{T}$$

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Conference graph adj mults: $S^2 = \dots = (2n - 2)I$

Conference Graph Doubling Theorem (Iverson, Jasper, M)

Sym conference mat order $n \implies$ ETF size $(n - 1) \times 2(n - 1)$

Ex: $q \equiv 1 \pmod{4}$ prime power \implies sym conference order $q + 1$
 \implies ETF size $q \times 2q$

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$d \times 2d$ ETF exists

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= previously unknown

Outline

- 1: Background
- 2: Doubling ETFs
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- 4: 2-circulant ETFs

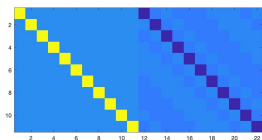
$d \times 2d$ ETF exists

1	2	3	4	5	6	7	8	9	10	<input type="checkbox"/> 11	12	13	14	15	16
<input type="checkbox"/> 17	18	19	20	21	22	23	24	25	26	27	28	<input type="checkbox"/> 29	30	31	32
33	34	<input type="checkbox"/> 35	36	37	38	<input type="checkbox"/> 39	40	41	42	<input type="checkbox"/> 43	44	45	46	<input type="checkbox"/> 47	48
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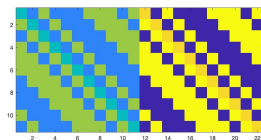
= previously unknown

Inspiration: Results of doubling

11×22

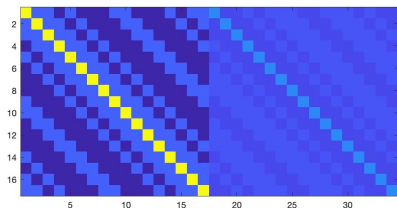


real part

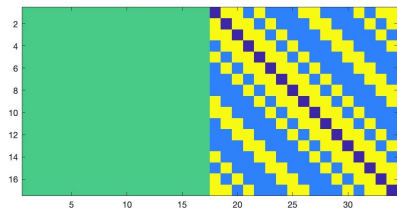


imaginary part

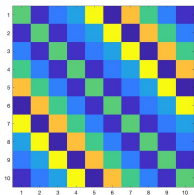
17×34



real part

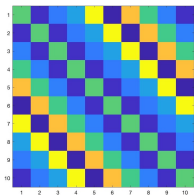


imaginary part



Circulant matrix:

- ▶ square
- ▶ rows cycle right
- ▶ columns cycle down

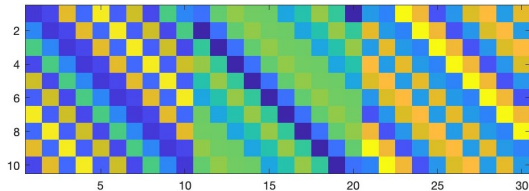


Circulant matrix:

- ▶ square
- ▶ rows cycle right
- ▶ columns cycle down

t -circulant matrix:

- ▶ $C = [C_1 \ \cdots \ C_t]$
- ▶ each C_j circulant
- ▶ $d \times td$



Other ETFs are t -circulant

Gerzon's bound

$$d \times n \text{ ETF over } \mathbb{F} \implies n \leq \begin{cases} \frac{d(d+1)}{2} & \mathbb{F} = \mathbb{R} \\ d^2 & \mathbb{F} = \mathbb{C} \end{cases}$$

Other ETFs are t -circulant

Gerzon's bound

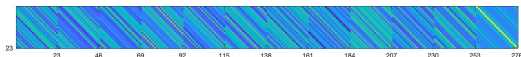
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3 × 6 over \mathbb{R}



7 × 28 over \mathbb{R}



23 × 276 over \mathbb{R}

Other ETFs are t -circulant

Gerzon's bound

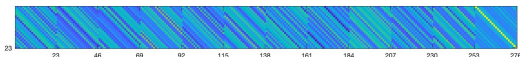
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3 × 6 over \mathbb{R}



7 × 28 over \mathbb{R}



23 × 276 over \mathbb{R}

~Zauner's conjecture

For every d , there exists a $d \times d^2$ ETF over \mathbb{C} that is d -circulant

Lemmens & Seidel, J. Algebra, 1973

Zauner, PhD thesis, 1999

Back to $d \times 2d$

$d \times 2d$ ETF from doubling is 2-circulant

1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
17	18	19	20	21	22	23	24	25	26	27	28	29	30	31	32
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= previously unknown

Back to $d \times 2d$

Medium $d \times 2d$ conjecture

For every d , there exists a $d \times 2d$ ETF that is 2-circulant

$d \times 2d$ ETF from doubling is 2-circulant

1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
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= previously unknown

Infinitely many 2-circulant $d \times 2d$ ETFs

Theorem (Iverson, Jasper, M)

q odd prime power \implies 2-circulant $d \times 2d$ ETF for $d = \begin{cases} \frac{1}{2}(q+1) \\ q+1 \end{cases}$

Infinitely many 2-circulant $d \times 2d$ ETFs

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Sketch:

- 1) Can detect 2-circulant from auto grp / signature mat

Infinitely many 2-circulant $d \times 2d$ ETFs

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Sketch:

- 1) Can detect 2-circulant from auto grp / signature mat
- 2) Apply for certain $d \times 2d$ ETFs:
 - a) $d = \frac{1}{2}(q+1)$ from Paley conference
 - b) $d = q+1$ from ETF doubling of (a)

So far

$d \times 2d$ ETF that is 2-circulant

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17	18	19	20	21	22	23	24	25	26	27	28	29	30	31	32
33	34	35	36	37	38	39	40	41	42	43	44	45	46	47	48
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81	82	83	84	85	86	87	88	89	90	91	92		94	95	96
97	98	99	100	101	102	103	104		106	107	108	109	110	111	112
113	114	115	116	117	118	119	120	121	122	123	124	125	126	127	128
129	130	131	132		134	135	136	137	138	139	140	141	142	143	144
145	146	147	148	149	150	151	152		154	155	156	157	158	159	160

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Want: 2-circulant $d \times 2d$ ETF $\Phi = [C \ D]$

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2-circulant structure reduces # angles to check:

$$\Phi\Phi^* = CC^* + DD^*, \quad \Phi^*\Phi = \begin{bmatrix} C^*C & C^*D \\ D^*C & D^*D \end{bmatrix}$$

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Strong $d \times 2d$ conjecture

For every $d \neq 4$, the set of 2-circulant $d \times 2d$ ETFs contains a real manifold of dimension $\lceil \frac{3d}{2} \rceil$

Theorem (Iverson, Jasper, M)

The strong $d \times 2d$ conjecture holds for all $d \leq 162$

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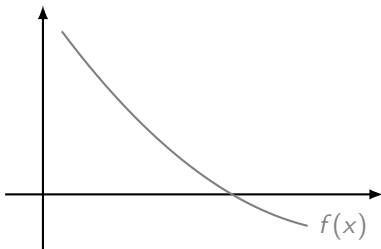
Proof: One d at a time, find approx soln \implies exact soln exists

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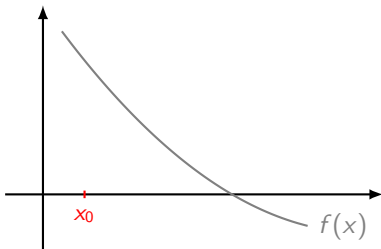


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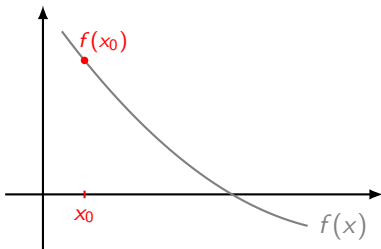


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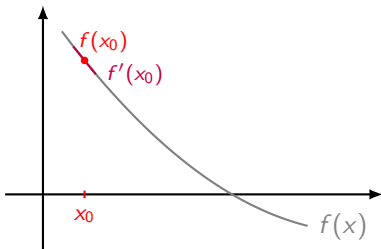


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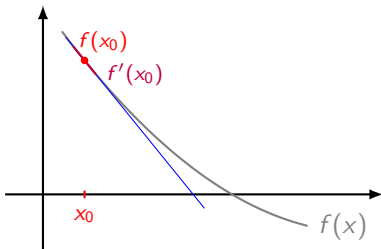


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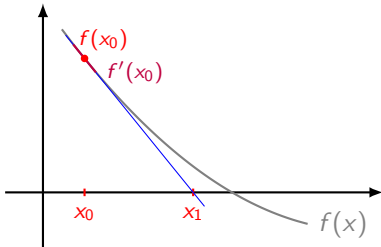


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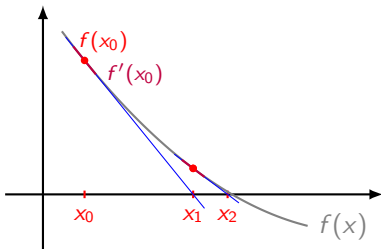


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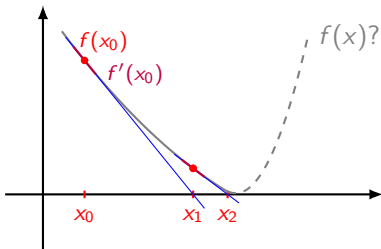


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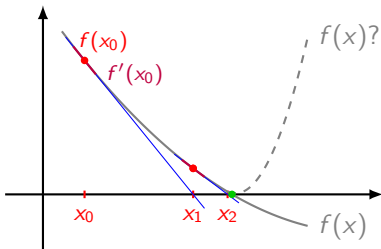


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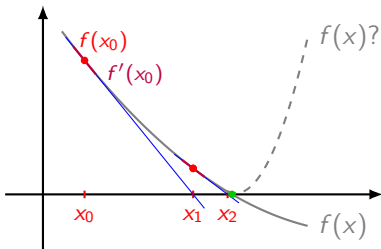
Control 2nd derivative \implies can't turn away!

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Furthermore: (local dim) = (# vars) - (# constraints)

$d \times 2d$ ETF that is 2-circulant

1	2	3	4	5	6	7	8	9	10	<input type="checkbox"/>	12	13	14	15	16
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Open questions

- ▶ Explicit solution to $d \times 2d$?
- ▶ How to leverage vars $>$ constraints to prove existence?
- ▶ Why is $d = 4$ different?
- ▶ Skew-Hadamard conjecture $\implies d \times 2d$ for $d \equiv 1 \pmod{4}$?
- ▶ Generic initializations converge to 2-circulant ETFs?

$d \times 2d$ ETF that is 2-circulant

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Questions?

On the optimal arrangement of $2d$ lines in \mathbb{C}^d

K. Fallon, J. W. Iverson

arXiv:2312.09975

More on the optimal arrangement of $2d$ lines in \mathbb{C}^d

J. W. Iverson, J. Jasper, D. G. Mixon

Coming soon!

Also, google **short fat matrices** for my research blog