

C&O 781 Topics in Quantum Information
Quantum Information Theory, Error-correction, and Cryptography
University of Waterloo
Spring 2008

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Assignment 3, Jun 27, 2008
Due: Jul 11, 2008

In the following questions, \mathcal{H} and \mathcal{K} are finite dimensional Hilbert spaces.

Question 1. [5 marks] For any $M \in L(\mathcal{H})$, prove that

$$\max_{\text{unitary } U \in L(\mathcal{H})} |\text{Tr}(UM)| = \|M\|_{\text{tr}}.$$

Conclude that

$$\|A + B\|_{\text{tr}} \leq \|A\|_{\text{tr}} + \|B\|_{\text{tr}}$$

for all $A, B \in L(\mathcal{H})$. (Therefore, the function $\|\cdot\|_{\text{tr}} : L(\mathcal{H}) \rightarrow \mathbb{R}$ is a norm.)

Question 2. [5 marks] Prove that for any density matrices $\rho_0, \rho_1 \in L(\mathcal{H})$,

$$\max_{\text{quantum states } \sigma \in L(\mathcal{H})} (\text{F}(\rho_0, \sigma) + \text{F}(\sigma, \rho_1)) = 1 + \sqrt{\text{F}(\rho_0, \rho_1)}.$$

Hint: first consider pure states ρ_0, ρ_1 .

Question 3. [5 marks] Show that for single qubit states $\rho_0, \rho_1 \in L(\mathcal{H})$,

$$1 - \text{F}(\rho_0, \rho_1) \leq \frac{1}{2} \|\rho_0 - \rho_1\|_{\text{tr}}.$$

Note that this is a stronger lower bound on trace distance in terms of fidelity than the general bound we saw in class.

Question 4. [5 marks] In the strong coin flipping protocol (with cheating probability $3/4$) we saw in class, we used qutrit states $|\psi_a\rangle \in \mathbb{C}^3 \otimes \mathbb{C}^3$. Find the best protocol of the same form, using bi-partite *qubit* states $|\phi_a\rangle$ given by:

$$|\phi_a\rangle = \cos \theta_a |00\rangle + \sin \theta_a |11\rangle,$$

as θ_a varies. (This in fact gives the best protocol with qubit states.)

Question 5. [5 marks] Suppose we replace the commitment states $|\psi_a\rangle$ in the strong coin-flipping protocol we saw in class by mixed states $\rho_a \in L(\mathbb{C}^d \otimes \mathbb{C}^d)$ of arbitrary dimension d . Explain what states and what final measurement by Bob would lead to a valid protocol. Then show that at least one part can cheat with probability at least $3/4$, regardless of what commitment state is used.