AN ALGORITHM FOR PACKING NON-ZERO A-PATHS IN GROUP-LABELED GRAPHS

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ABSTRACT. Let $G = (V, E)$ be an oriented graph whose edges are labeled by the elements of a group Γ and let $A \subseteq V$. An *A-path* is a path whose ends are both in A. The *weight* of a path P in G is the sum of the group values on forward oriented arcs minus the sum of the backward oriented arcs in P. (If Γ is not abelian, we sum the labels in their order along the path.) We give an efficient algorithm for finding a maximum collection of vertex-disjoint A-paths each of non-zero weight. When $A = V$ this problem is equivalent to the maximum matching problem.

1. INTRODUCTION

Let Γ be a group; we will use additive notation for groups, although they need not be abelian. A *Γ-labeled graph* is a graph G in which each edge $e = uv \in E(G)$ is assigned weights $\omega_G(e, u), \omega_G(e, v) \in \Gamma$ where $\omega_G(e, u) = -\omega_G(e, v)$. Let G be a Γ-labeled graph and let $A \subseteq V(G)$. An *A-path* is a path, with at least one edge, whose ends are both in A. Now, if $P = (v_0, e_1, v_1, e_2, v_2, \ldots, e_k, v_k)$ is a path in G, then the *weight of P*, denoted $\omega_G(P)$, is defined to be $\sum_{i=1}^{k} \omega_G(e_i, v_i)$.

We are interested in the maximum number of vertex-disjoint A-paths each of non-zero weight, which we denote by $\nu(G, A)$. Chudnovsky *et al.* [1] gave a min-max theorem for $\nu(G, A)$; they also discuss motivation for the non-zero A-paths problem. In particular, they show that Mader's $\mathcal{S}$-path problem [4] is a special case. The only previously known algorithm for Mader's $\mathcal{S}$-path problem was obtained by Lovász via a reduction to linear matroid matching [2]. We present an algorithm for finding a maximum collection of vertex-disjoint non-zero A-paths that runs in time $O(|V(G)|^6)$. In our complexity calculations, group

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operations (such as addition and comparison) are treated as elementary operations. Our algorithm is similar to an algorithm of Lovász and Plummer [3, p376] for finding a maximum matching. Lovász and Plummer cleverly abstract an algorithm from what would otherwise appear to be a nonconstructive proof of the Edmonds-Gallai Structure Theorem (see [3]). Using a similar approach, we obtain an algorithm from our proof of Theorem 1.3, which is a structure theorem for non-zero A-paths. Theorem 1.3 is closely related to a structure theorem of Sebő and Szegő [5] for Mader's $\mathcal{S}$-path problem; our results were, however, obtained independently.

Let $E_0(G, A)$ denote the set of all edges $e = uv \in E$ whose ends are both in A and that have $\omega_G(e, v) = 0$; note that deleting such edges does not affect ν. Let $\mathrm{def}(G, A) = |A| - 2\nu(G, A)$; we call this the *deficiency*. Let $\mathrm{odd}(G, A)$ denote the number of components H of $G - E_0(G, A)$ with $|V(H) \cap A|$ odd. Finally let $X, A' \subseteq V(G)$ such that $A \cup X \subseteq A'$. It is straightforward to see that

$$\begin{aligned} \mathrm{def}(G, A) &\geq \mathrm{def}(G, A') - |A' - A| \\ &\geq \mathrm{def}(G - X, A' - X) - |A' - A| - |X| \\ &\geq \mathrm{odd}(G - X, A' - X) - |A' - A| - |X|. \end{aligned}$$

Let $x \in V$ and let $\delta \in \Gamma$. We will construct a new Γ-labeled graph G' from G by changing the labels as follows. For each edge $e = uv$ in G we define

$$\omega_{G'}(e, u) = \begin{cases} \omega_G(e, u) + \delta, & \text{if } u = x \\ -\delta + \omega_G(e, u), & \text{if } v = x \\ \omega_G(e, u), & \text{otherwise.} \end{cases}$$

We say that G' is obtained from G by *shifting* by δ at x. Note that, if $x \notin A$, then this shift does not change the weight of any A-path (even when Γ is non-abelian). If G' is a Γ-labeled graph obtained by a sequence of shifting operations on vertices not in A, then we say that G and G' are *A-equivalent*. The main theorem in [1] is:

Theorem 1.1. *Let Γ be a group, let G be a Γ-labeled graph, and let $A \subseteq X$. Then there exists a Γ-labeled graph G' that is A-equivalent to G and there exist sets $X, A' \subseteq V$ with $A \cup X \subseteq A'$ such that*

$$\mathrm{def}(G, A) = \mathrm{odd}(G' - X, A' - X) - |A' - A| - |X|.$$

Our structure theorem provides a canonical choice for A' and X in Theorem 1.1. Before stating the structure theorem we need some definitions; we start by clarifying what we mean by a path.

A *path* is a sequence $P = (v_0, e_1, v_1, e_2, v_2, \ldots, e_k, v_k)$ where $v_0, \ldots, v_k$ are distinct vertices of G and e_i has ends v_{i-1} and v_i for each $i \in$

$\{1, \dots, k\}$. Thus P is ordered in that it has distinguished *start* (v_0) and *end* (v_k). The path $(v_k, e_k, v_{k-1}, \dots, v_1, e_1, v_0)$ is denoted by $\bar{P}$. We allow paths consisting of a single vertex; we refer to such paths as *trivial*.

An A-*collection* is a set Π of vertex disjoint paths such that:

1. each vertex in A is either the start or the end of a path in Π,
2. the start of each path $P \in \Pi$ is in A, and
3. if $P \in \Pi$ is non-trivial and has its end in A, then $\omega_G(P) \neq 0$.

A path $P \in \Pi$ is *loose* if it is trivial or its end is not in A; thus each path in Π is either an A-path or it is loose (not both). The *value* of an A-collection Π, denoted $\mathrm{val}_A(\Pi)$ or $\mathrm{val}(\Pi)$, is the number of A-paths that it contains. An A-collection is *optimal* if $\mathrm{val}(\Pi) = \nu(G, A)$; note that there are optimal A-collections. Let $\mathcal{P}(G, A)$ denote the set of all A-collections and let $\mathcal{P}^*(G, A)$ denote the set of all optimal A-collections.

Let Π be an A-collection and let $B_A(\Pi)$ (or $B(\Pi)$) denote the set of pairs $(v, \omega_G(P))$ where v is the end of a loose path $P \in \Pi$. Note that $|B(\Pi)| = |A| - 2\mathrm{val}_A(\Pi)$. Now let $\mathcal{R}(G, A) = \cup(B(\Pi) \,:\, \Pi \in \mathcal{P}^*(G, A))$; the pairs in $\mathcal{R}(G, A)$ are called *reachable pairs*.

For each vertex $v \in V(G)$, we let $\Gamma(G, A, v) = \{\alpha \in \Gamma \,:\, (v, \alpha) \in \mathcal{R}(G, A)\}$. Now we let

$$\begin{aligned} D_1(G, A) &= \{v \in V(G) \,:\, |\Gamma(G, A, v)| = 1\}, \\ D_2(G, A) &= \{v \in V(G) \,:\, |\Gamma(G, A, v)| \geq 2\}, \text{ and} \\ D(G, A) &= D_1(G, A) \cup D_2(G, A); \end{aligned}$$

$D(G, A)$ is the set of *reachable* vertices. Note that $D_1(G, A)$ and $D_2(G, A)$ are not affected by shifting on a vertex $v \notin A$.

To make use of the coming structure theorem, we need the following easy lemma.

Lemma 1.2. *Let G be a Γ-labeled graph and let $A \subseteq V(G)$. Then there exists a Γ-labeled graph G' that is A-equivalent to G and such that:*

(1) *for each $v \in D_1(G', A)$, $\Gamma(G', A, v) = \{0\}$, and*
(2) *for each $u \in N_{G'}(D(G', A)) - A$, there exists $uv = e \in E(G')$ such that $\omega_{G'}(e, v) \in \Gamma(G', A, v)$.*

Proof. Suppose that $v \in D_1(G, A)$ and that $\Gamma(G, A, v) = \{\alpha\}$. If $v \in A$, then $\alpha = 0$. On the other hand, if $v \notin A$ and G' is obtained from G by shifting by α at v, then $\Gamma(G', A, v) = \{0\}$ and $\Gamma(G', A, y) = \Gamma(G, A, y)$ for all $y \in V(G) - \{v\}$.

Now suppose that $uv = e \in E(G)$ where $u \notin A \cup D(G, A)$ and $v \in D(G, A)$. Let $\alpha \in \Gamma(G, A, v)$ and let G' be the Γ-labeled graph obtained from G by shifting by $\omega_G(e,v)-\alpha$ at u. Then $\omega_{G'}(e,v)=\alpha$ and $\Gamma(G',A,y)=\Gamma(G,A,y)$ for all $y\in V(G)$. □

We can now state our structure theorem.

Theorem 1.3. *Let Γ be a group, let G be a Γ-labeled graph, and let $A\subseteq V(G)$. Now let $A'=A\cup N_G(D(G,A))\cup D_1(G,A)$ and let $X=N_{G-E_0(G,A')}(D(G,A))$. If (G,A) satisfies:*

1. *for each $v\in D_1(G,A)$, $\Gamma(G,A,v)=\{0\}$, and*
2. *for each $u\in N_G(D(G,A))-A$, there exists $uv=e\in E(G)$ such that $\omega_G(e,v)\in\Gamma(G,A,v)$,*

then $\mathrm{def}(G,A)=\mathrm{odd}(G-X,A'-X)-|A'-A|-|X|$.

2. Proof of the structure theorem

In this section we outline a proof of the structure theorem; this outline is intended to motivate the main steps in the algorithm. Throughout the rest of the paper we let Γ be a group, we let G be a group labeled graph, and we let $A\subseteq V(G)$.

We begin with two easy observations:

2.1. *If $u\in A-D(G,A)$, then $\nu(G-u,A-\{u\})=\nu(G,A)-1$ and $\mathcal{R}(G,A)\subseteq\mathcal{R}(G-u,A-\{u\})$.*

2.2. *If $u\in V(G)-A$ and $\Gamma(G,A,u)\subseteq\{0\}$, then $\nu(G,A\cup\{u\})=\nu(G,A)$ and $\mathcal{R}(G,A)\subseteq\mathcal{R}(G,A\cup\{u\})$.*

We are interested in elements for which the inclusions, in 2.1 and 2.2, hold with equality; the next two results identify some such elements. We will not prove these lemmas now since they follow immediately from more general results (Lemmas 4.3 and 4.4) proved later.

Lemma 2.3. *Let $u\in A-D(G,A)$. If there exists $uv=e\in E(G)$ and $\alpha\in\Gamma(G,A,v)$ such that $\omega_G(e,u)\neq-\alpha$, then $\nu(G-u,A-\{u\})=\nu(G,A)-1$ and $\mathcal{R}(G,A)=\mathcal{R}(G-u,A-\{u\})$.*

Lemma 2.4. *Let $u\in V(G)-A$ where $\Gamma(G,A,u)\subseteq\{0\}$. If there exists $uv=e\in E(G)$ and $\alpha\in\Gamma(G,A,v)$ such that $\omega_G(e,u)=-\alpha$, then $\nu(G,A\cup\{u\})=\nu(G,A)$ and $\mathcal{R}(G,A)=\mathcal{R}(G,A\cup\{u\})$.*

With the main ingredients in place, we can begin the proof of the structure theorem. Suppose that:

1. for each $v\in D_1(G,A)$, $\Gamma(G,A,v)=\{0\}$, and
2. for each $u\in N_G(D(G,A))-A$, there exists $uv=e\in E(G)$ such that $\omega_G(e,v)\in\Gamma(G,A,v)$.

Now let $A' = A \cup N_G(D(G,A)) \cup D_1(G,A)$ and $X = N_{G-E_0(G,A')}(D(G,A))$.

Lemma 2.5. *$\nu(G,A) = \nu(G-X, A'-X) + |X|$ and $\mathcal{R}(G,A) = \mathcal{R}(G-X, A')$. Hence* $\mathrm{def}(G,A) = \mathrm{def}(G-X, A'-X) - |A'-A| - |X|$.

Proof. First we consider $u \in N_G(D(G,A)) - A$. By (2), there exists $uv = e \in E(G)$ such that $\omega_G(e,v) \in \Gamma(G,A,v)$. Then, by Lemma 2.4, $\nu(G, A\cup\{u\}) = \nu(G,A)$ and $\mathcal{R}(G,A) = \mathcal{R}(G, A\cup\{u\})$.

Now consider $u \in D_1(G,A) - A$. By (1), we have $\Gamma(G,A,u) = \{0\}$. Thus there exists $\Pi \in \mathcal{P}^*(G,A)$ such that $(u,0) \in B(\Pi)$. Let $P \in \Pi$ be the path containing u, let v be the vertex preceding u on P, let P_v be the initial subpath of P ending at v, and let $e = uv$. Note that $\omega_G(P_v) - \omega_G(e,v) = \omega_G(P) = 0$. Thus $\omega_G(P_v) = \omega_G(e,v)$. Let $\Pi_v = (\Pi - \{P\}) \cup \{P_v\}$. Now Π_v is an optimal A-collection with $(v, \omega_G(P_v)) \in B(\Pi_v)$. Therefore $\omega_G(e,v) \in \Gamma(G,A,v)$. Hence, by Lemma 2.4, $\nu(G, A\cup\{u\}) = \nu(G,A)$ and $\mathcal{R}(G,A) = \mathcal{R}(G, A\cup\{u\})$.

This proves that $\nu(G,A) = \nu(G,A')$ and $\mathcal{R}(G,A) = \mathcal{R}(G,A')$. Moreover, note that $X \subseteq A'$. Now consider $u \in X$. By definition, there exists $uv = e \in E(G) - E_0(G,A')$ where $v \in D(G,A')$. Therefore there exists $\alpha \in \Gamma(G,A',v)$ such that $\omega_G(e,v) \neq \alpha$. Therefore, by Lemma 2.3, $\nu(G-u, A'-\{u\}) = \nu(G,A') - 1$ and $\mathcal{R}(G,A') = \mathcal{R}(G-u, A'-\{u\})$. It follows that $\nu(G,A) = \nu(G-X, A'-X) + |X|$ and $\mathcal{R}(G,A) = \mathcal{R}(G-X, A')$, as required. □

We need one more definition. A *critical pair* (G,A) consists of a Γ-labeled graph G and a set $A \subseteq V(G)$ such that G is connected, $D_1(G,A) = A$, $D_2(G,A) = V(G) - A$, and $E_0(G,A) = \emptyset$.

Now let $A_1 = A' - X$ and $G_1 = G - X$. The next lemma follows from the definition of G_1 and A_1.

Lemma 2.6. *For each component H of $G_1 - E_0(G_1, A_1)$, either* $\mathrm{def}(H, V(H)\cap A_1) = 0$ *or $(H, V(H)\cap A_1)$ is critical.*

Proof. Let $G_2 = G_1 - E_0(G_1, A_1)$. By Lemma 2.5, $D(G_2, A_1) = D(G,A)$. Note that, if $uv = e \in E(G)$ with $u \in V(G) - D(G,A)$ and $v \in D(G,A)$, then either $u \in X$ or $e \in E_0(G,A')$. Thus, if H is a component of G_2, then either $V(H) \subseteq D(G_2,A_1)$ or $V(H) \cap D(G_2,A_1) = \emptyset$. If $V(H) \cap D(G_2,A_1) = \emptyset$, then $\mathrm{def}(H, V(H)\cap A_1) = 0$. Thus we may assume that $V(H) \subseteq D(G_2,A_1)$. Note that, since H is a component of G_2, $D_1(H, V(H)\cap A_1) = D_1(G_2,A_1) \cap V(H)$ and $D_2(H, V(H)\cap A_1) = D_2(G_2,A_1)\cap V(H)$. By the definition of A', a vertex $v \in D(G_2,A_1)$ is in $D_1(G_2,A_1)$ if and only if $v \in A_1$. Hence $(H, V(H)\cap A_1)$ is critical, as required. □

The final lemma was proved in [1]; we prove a more general lemma later (see 4.5).

Lemma 2.7. *If* (G, A) *is a critical pair, then* $\mathrm{def}(G, A) = 1$ *and, hence,* $|A|$ *is odd.*

It follows from Lemmas 2.6 and 2.7 that $\mathrm{def}(G_1, A_1) = \mathrm{odd}(G_1, A_1)$. Therefore

$$\mathrm{def}(G, A) = \mathrm{odd}(G - X, A' - X) - |A' - A| - |X|.$$

This completes the proof of the structure theorem.

3. The exchange property

Chudnovsky *et al.* [1] proved that $\{B(\Pi) : \Pi \in \mathcal{P}^*(G, A)\}$ is the set of bases of a matroid. The following lemma extends that result by providing an exchange property on all A-collections. The proof is essentially the same as the proof given in [1]. (For sets A and B, we let $A \Delta B = (A - B) \cup (B - A)$.)

Lemma 3.1. *Let* $\Pi_1, \Pi_2 \in \mathcal{P}(G, A)$ *and let* $p_1 \in B(\Pi_1) - B(\Pi_2)$. *Then there exists* $\Pi_1' \in \mathcal{P}(G, A)$ *and* $p_2 \in B(\Pi_1) \cup B(\Pi_2)$ *such that* $B(\Pi_1') = B(\Pi_1)\Delta\{p_1, p_2\}$. *Moreover, given* Π_1, Π_2, *and* p_1, *we can find* Π_1' *and* p_2 *in* $O(|V(G)|^2)$ *time.*

Proof. Suppose, by way of contradiction, that there exist

3.1.1. $\Pi_1, \Pi_2 \in \mathcal{P}(G, A)$ *and* $p_1 = (u, \alpha) \in B(\Pi_1) - B(\Pi_2)$ *such that there do not exist* $\Pi_1' \in \mathcal{P}(G, A)$ *and* $p_2 \in B(\Pi_1) \cup B(\Pi_2)$ *with* $B(\Pi_1') = B(\Pi_1)\Delta\{p_1, p_2\}$.

3.1.2. *We choose* Π_1, Π_2, *and* (u, α) *satisfying 3.1.1 with* $|E(\Pi_1) \cup E(\Pi_2)|$ *as small as possible.*

We use the following claim repeatedly.

3.1.3. *There do not exist* $\Pi_1' \in \mathcal{P}(G, A)$ *and* $p_2 \in V(G) \times \Gamma$ *such that* $B(\Pi_1') = B(\Pi_1)\Delta\{p_1, p_2\}$ *and* $|E(\Pi_1') \cup E(\Pi_2)| < |E(\Pi_1) \cup E(\Pi_2)|$.

Proof of claim. Suppose otherwise. By 3.1.1, $p_2 \notin B(\Pi_1) \cup B(\Pi_2)$. However, $|E(\Pi_1') - E(\Pi_2)| < |E(\Pi_1) - E(\Pi_2)|$. So, by 3.1.2, Π_1', Π_2, and p_2 do not satisfy 3.1.1. That is, there exists an element $p_3 \in B(\Pi_2) - B(\Pi_1')$ such that $B(\Pi_1')\Delta\{p_2, p_3\} \in \mathcal{B}$. However, $B(\Pi_1)\Delta\{p_1, p_3\} = B(\Pi_1')\Delta\{p_2, p_3\} \in \mathcal{B}$, contradicting 3.1.1. □

Let $p_1 = (u, \alpha)$ and let $P = (v_0, e_1, v_1, \ldots, e_k, v_k)$ be the path in Π_1 ending at u. By possibly reversing the order, we may assume that there is a path $P' = (v'_0, e'_1, v'_1, \ldots, e'_l, v'_l)$ in Π_2 that starts at v_0. Suppose that P is not contained in P'. Now let Π'_1 be the A-collection obtained from Π_1 by replacing P with the trivial path (v_0). Note that $B(\Pi'_1) = B(\Pi_1)\Delta\{p_1, (v_0, 0)\}$ and $|E(\Pi'_1) \cup E(\Pi_2)| < |E(\Pi_1) \cup E(\Pi_2)|$, contradicting 3.1.3. Hence P is contained in P'.

Suppose that P' is disjoint from each path in $\Pi_1 - \{P\}$ and let Π'_1 be obtained from Π_1 by replacing P with P'. Note that $B(\Pi'_1) = B(\Pi_1)\Delta\{p_1, (v'_l, \omega_G(P'))\}$ and $(v'_l, \omega_G(P')) \in B(\Pi_2)$, contradicting 3.1.1. Therefore there is some vertex that is both on P' and on a path in Π_1 other than P; let v'_i be the first such vertex on P' and let $Q = (u_0, f_1, u_1, \ldots, f_m, u_m)$ be the path of Π_1 containing v'_i. Suppose that $u_j = v'_i$. We consider two cases.

Case 1: *Q is a loose path.*

Let P_1 be the A-path contained in $G[E(P') \cup E(Q)]$ and let P_2 be the path in $G[E(P') \cup E(Q)]$ that starts at v_0 and ends at u_m.

Case 1.1: $\omega_G(P_1) \neq 0$.

Let $\Pi'_1 = (\Pi_1 - \{P, Q\}) \cup P_1$. Note that $B(\Pi'_1) = B(\Pi_1) - \{p_1, (u_m, \omega_G(Q))\}$ and $(u_m, \omega_G(Q)) \in B(\Pi_1)$, contradicting 3.1.1.

Case 1.2: $\omega_G(P_1) = 0$.

Thus $\omega(P'[v'_0, v'_i]) = \omega(Q[u_0, u_j])$ and, hence, $\omega(P_2) = \omega(Q)$. Now let Π'_1 be the A-collection obtained from Π_1 by replacing P and Q with P_2 and the trivial path (u_0). Note that $B(\Pi'_1) = (B(\Pi_1) - \{p_1\}) \cup \{(u_0, 0)\}$. Moreover, since $\omega_G(P_1) = 0$, $P_1 \neq P'$. Thus there is an edge of $Q[u_0, u_j]$ that is not in $E(\Pi_2)$. So, $|E(\Pi'_1) \cup E(\Pi_2)| < |E(\Pi_1) \cup E(\Pi_2)|$, contradicting 3.1.3.

Case 2: *Q is an A-path.*

Let P_1 and P_2 be the A-paths in $G[E(P') \cup E(Q)]$ that both start at v_0 and that end with u_0 and u_m respectively. Note that $\omega(P_1) + \omega(Q) + \omega(\bar{P}_2) = 0$ and $\omega(Q) \neq 0$, so either $\omega(P_1) \neq 0$ or $\omega(P_2) \neq 0$. Moreover, either P' is loose (and hence different from P_1 and P_2) or $\omega(P') \neq 0$. Thus either $\omega(P_1) \neq 0$ and $P_2 \neq P'$ or $\omega(P_2) \neq 0$ and $P_1 \neq P'$. By possibly swapping P_1 and P_2 and reversing Q, we may assume that $\omega(P_2) \neq 0$ and $P_1 \neq P'$. Let Π'_1 be the A-collection obtained from Π_1 by replacing P and Q with P_2 and the trivial path (u_0). Note that $B(\Pi'_1) = (B(\Pi_1) - \{p_1\}) \cup \{(u_0, 0)\}$. Moreover, since $P_1 \neq P'$ there is an edge of $Q[u_0, u_j]$ that is not in $E(\Pi'_1) \cup E(\Pi_2)$. Thus $|E(\Pi'_1) \cup E(\Pi_2)| < |E(\Pi_1) \cup E(\Pi_2)|$, contradicting 3.1.3. This final contradiction completes the proof.

The above proof can easily be made algorithmic with the desired running time. □

We now prove a useful application of the exchange property.

Lemma 3.2. *Let* $\Pi_1, \Pi_2 \in \mathcal{P}(G, A)$ *and let* $B_1 \subseteq B_A(\Pi_1)$*. Then there exists* $\Pi_3 \in \mathcal{P}(G, A)$ *such that either:*

(1) $\mathrm{val}(\Pi_3) = \mathrm{val}(\Pi_1)$ *and* $B_1 \subseteq B(\Pi_3)$ *and* $B(\Pi_3) - B_1 \subseteq B(\Pi_2)$*, or*

(2) $\mathrm{val}(\Pi_3) = \mathrm{val}(\Pi_1) + 1$ *and* $|B(\Pi_3) \cap B_1| \geq |B_1| - 1$*.*

Moreover, we can find such Π_3 *in* $O(|V(G)|^3)$ *time.*

Proof. We assume that:

($*$) *Among all* $\Pi_1 \in \mathcal{P}(G, A)$ *with* $B_1 \subseteq B(\Pi_1)$ *we choose* Π_1 *minimizing* $|B(\Pi_1) - B(\Pi_2)|$.

We may assume that there exists $p_1 \in B(\Pi_1) - (B_1 \cup B(\Pi_2))$, since otherwise $\Pi_3 := \Pi_1$ satisfies (1). By the exchange property, there exists $\Pi_3 \in \mathcal{P}(G, A)$ and $p_2 \in B(\Pi_1) \cup B(\Pi_2)$ such that $B(\Pi_3) = B(\Pi_1)\Delta\{p_1, p_2\}$.

Case 1: $p_2 \in B(\Pi_1)$.

Thus $\mathrm{val}(\Pi_3) = \mathrm{val}(\Pi_1) + 1$ and $|B(\Pi_3) \cap B_1| \geq |B_1| - 1$, satisfying (2).

Case 2: $p_2 \in B(\Pi_2) - B(\Pi_1)$.

Thus $\mathrm{val}(\Pi_3) = \mathrm{val}(\Pi_1)$, $B_1 \subseteq B(\Pi_3)$, and $|B(\Pi_3) - B(\Pi_2)| < |B(\Pi_1) - B(\Pi_2)|$, contradicting ($*$).

That completes the proof; this proof can clearly be made algorithmic with the desired running time. □

The following two lemmas are consequences of Lemma 3.2.

Lemma 3.3. *Let* $\Pi_1, \Pi_2 \in \mathcal{P}(G, A)$ *with* $\mathrm{val}(\Pi_2) = \mathrm{val}(\Pi_1) + 1$*, let* $uv = e \in E(G)$*, let* (u, α) *and* p *be distinct elements of* $B(\Pi_1)$*, and let* $(v, \beta) \in B(\Pi_2)$ *where* $\alpha + \omega_G(e, v) - \beta \neq 0$*. Then there exists* $\Pi_3 \in \mathcal{P}(G, A)$ *such that* $\mathrm{val}(\Pi_3) = \mathrm{val}(\Pi_2)$ *and either* $(u, \alpha) \in B(\Pi_3)$ *or* $p \in B(\Pi_3)$*. Moreover, we can find such* Π_3 *in* $O(|V(G)|^3)$ *time.*

Proof. By Lemma 3.2 with $B_1 = \{p, (u, \alpha)\}$, we get one of the following two cases.

Case 1: *There exists* $\Pi \in \mathcal{P}(G, A)$ *such that* $\mathrm{val}(\Pi) = \mathrm{val}(\Pi_1)$ *and* $B_1 \subseteq B(\Pi)$ *and* $B(\Pi) - B_1 \subseteq B(\Pi_2)$*.*

Thus $p, (u, \alpha), (v, \beta) \in B(\Pi)$. Let P_u and P_v be the loose paths in Π ending at u and v respectively. Now let $P = (P_1, e, \bar{P}_2)$. Note that P is an A-path and $\omega_G(P) = \alpha + \omega_G(e, v) - \beta \neq 0$. Now let $\Pi_3 = (\Pi - \{P_u, P_v\}) \cup \{P\}$. Note that $\mathrm{val}_A(\Pi_3) = \mathrm{val}(\Pi_2)$ and $p \in B(\Pi_3)$, as required.

Case 2: *There exists* $\Pi_3 \in \mathcal{P}(G, A)$ *such that* $\text{val}(\Pi_3) = \text{val}(\Pi_2)$ *and* $|B(\Pi_3) \cap B_1| \geq |B_1| - 1$.

Thus either $(u, \alpha) \in B(\Pi_3)$ or $p \in B(\Pi_3)$, as required.

This proof is clearly constructive with the desired running time. □

The next lemma is a direct consequence of Lemma 3.2; we omit the easy proof.

Lemma 3.4. *Let* $\Pi_1, \Pi_2 \in \mathcal{P}(G, A)$ *with* $\text{val}(\Pi_2) = \text{val}(\Pi_1)$*, let* $p_1 \in B(\Pi_1)$*, and let* p_2 *and* p_3 *be distinct elements of* $B(\Pi_2)$*. Then there exists* $\Pi_3 \in \mathcal{P}(G, A)$ *such that either:*

(1) $\text{val}(\Pi_3) = \text{val}(\Pi_1)$*,* $p_1 \in B(\Pi_3)$*, and either* $p_2 \in B(\Pi_3)$ *or* $p_3 \in B(\Pi_3)$*, or*

(2) $\text{val}(\Pi_3) = \text{val}(\Pi_1) + 1$.

Moreover, we can find such Π_3 *in* $O(|V(G)|^3)$ *time.*

4. Key lemmas

In this section we prove constructive analogues of some of the lemmas in Section 2.

Throughout this section we let G be a Γ-labeled graph, $A \subseteq V(G)$, and $\mathcal{P} \subseteq \mathcal{P}(G, A)$. We use the following definitions:

$$\begin{aligned}
\nu(\mathcal{P}, A) &= \max(\text{val}_A(\Pi) \,:\, \Pi \in \mathcal{P}),\\
\text{def}(\mathcal{P}, A) &= |A| - 2\nu(\mathcal{P}, A),\\
\mathcal{P}^* &= \{\Pi \in \mathcal{P} \,:\, \text{val}_A(\Pi) = \nu(\mathcal{P}, A)\}, \text{ and}\\
\mathcal{R}(\mathcal{P}, A) &= \cup(B_A(\Pi) \,:\, \Pi \in \mathcal{P}^*).
\end{aligned}$$

Now, for each $v \in V(G)$, we let

$$\Gamma(\mathcal{P}, A, v) = \{\gamma \in \Gamma \,:\, (v, \gamma) \in \mathcal{R}(\mathcal{P}, A)\}.$$

In addition, we define:

$$\begin{aligned}
D_1(\mathcal{P}, A) &= \{v \in V(G) \,:\, |\Gamma(\mathcal{P}, A, v)| = 1\},\\
D_2(\mathcal{P}, A) &= \{v \in V(G) \,:\, |\Gamma(\mathcal{P}, A, v)| > 1\}, \text{ and}\\
D(\mathcal{P}, A) &= D_1(\mathcal{P}, A) \cup D_2(\mathcal{P}, A).
\end{aligned}$$

We begin with some easy observations relating to 2.1 and 2.2:

4.1. *Let* $u \in A - D(\mathcal{P}, A)$*. If there exists* $\Pi \in \mathcal{P}(G - u, A - \{u\})$ *such that* $\text{val}_{A-\{u\}}(\Pi) = \nu(\mathcal{P}, A)$*, then there exists* $\Pi' \in \mathcal{P}(G, A)$ *such that* $(u, 0) \in \mathcal{R}(\mathcal{P} \cup \{\Pi'\}, A)$.

4.2. *Let $u \in V(G) - A$ with $\Gamma(\mathcal{P}, A, u) \subseteq \{0\}$. If there exists $\Pi \in \mathcal{P}(G, A \cup \{u\})$ such that $\mathrm{val}_{A\cup\{u\}}(\Pi) = \nu(\mathcal{P}, A) + 1$, then there exists $\Pi' \in \mathcal{P}(G, A)$ such that either $\mathrm{val}_A(\Pi') > \nu(\mathcal{P}, A)$ or $\mathrm{val}_A(\Pi') = \nu(\mathcal{P}, A)$ and there exists $\alpha \in \Gamma - \{0\}$ such that $(u, \alpha) \in \mathcal{R}(\mathcal{P} \cup \{\Pi'\}, A)$.*

The next result generalizes Lemma 2.3.

Lemma 4.3. *Let $u \in A - D(\mathcal{P}, A)$, $uv = e \in E(G)$, and $\alpha \in \Gamma(\mathcal{P}, A, v)$ such that $\omega_G(e, u) \neq -\alpha$. If there exists $\Pi \in \mathcal{P}(G-u, A-\{u\})$ such that $\mathrm{val}_{A-\{u\}}(\Pi) = \nu(\mathcal{P}, A) - 1$ and there exists $p \in B_{A-\{u\}}(\Pi) - \mathcal{R}(\mathcal{P}, A)$, then there exists $\Pi' \in \mathcal{P}(G, A)$ such that $\mathrm{val}_A(\Pi') = \nu(\mathcal{P}, A)$ and either $(u, \alpha + \omega_G(e, u)) \in B(\Pi')$ or $p \in B(\Pi')$. Moreover, if $|\mathcal{P}| \leq 2|V(G)|$, then we can find such Π' in $O(|V(G)|^3)$ time.*

Proof. Let Π_1 be the A-collection obtained by adding the trivial path (u) to Π. Note that $\mathrm{val}_A(\Pi_1) = \nu(\mathcal{P}, A) - 1$ and $p, (u, 0) \in B_A(\Pi_1)$. Let $\Pi_2 \in \mathcal{P}^*$ with $(v, \alpha) \in B_A(\Pi_2)$. Now $\mathrm{val}_A(\Pi_2) = \mathrm{val}_A(\Pi_1) + 1$. Therefore, by Lemma 3.3, we find $\Pi' \in \mathcal{P}(G, A)$ with $\mathrm{val}_A(\Pi') = \nu(\mathcal{P}, A)$ and either $(u, 0) \in B_A(\Pi')$ or $p \in B_A(\Pi')$. Now Π' satisfies the lemma.

This proof is clearly constructive with the desired running time. □

The next result generalizes Lemma 2.4.

Lemma 4.4. *Let $u \in V(G) - A$ with $\Gamma(\mathcal{P}, A, u) \subseteq \{0\}$, let $uv = e \in E(G)$ with $\omega_G(e, v) \in \Gamma(\mathcal{P}, A, v)$. If there exists $\Pi \in \mathcal{P}(G, A \cup \{u\})$ such that $\mathrm{val}_{A\cup\{u\}}(\Pi) = \nu(\mathcal{P}, A)$ and there exists $p \in B_{A\cup\{u\}}(\Pi) - \mathcal{R}(\mathcal{P}, A)$, then there exists $\Pi' \in \mathcal{P}(G, A)$ such that either $\mathrm{val}_A(\Pi') > \nu(\mathcal{P}, A)$ or $\mathrm{val}_A(\Pi') = \nu(\mathcal{P}, A)$ and either $p \in B(\Pi')$ or there exists $(u, \alpha) \in B(\Pi')$ with $\alpha \neq 0$. Moreover, if $|\mathcal{P}| \leq 2|V(G)|$, then we can find such Π' in $O(|V(G)|^3)$ time.*

Proof. Note that, if $v \in A$, then, since $\omega_G(e, v) \in \Gamma(\mathcal{P}, A, v)$, we have $\omega_G(e, v) = 0$. On the other hand, if $v \notin A$, then, by possibly shifting, we may assume that $\omega_G(e, v) = 0$. Let $p = (w, \delta)$. We break the proof into the following cases.

Case 1: *There exists $\Pi_1 \in \mathcal{P}(G, A \cup \{u\})$ with $\mathrm{val}_{A\cup\{u\}}(\Pi_1) = \nu(\mathcal{P}, A)$ and $p \in B_{A\cup\{u\}}(\Pi_1)$, such that u is not the start of the loose path in Π_1 containing w.*

There is a path $P \in \Pi_1$ whose start or end is u. Suppose that P is a loose path with respect to $A \cup \{u\}$; thus u is the start of P and P does not contain w. Then $\Pi' := \Pi_1 - \{P\}$ satisfies the lemma. Therefore we may assume that P is an $A \cup \{u\}$-path; furthermore, by possibly reversing P, we may assume that u is the end of P. Let $\alpha = \omega_G(P)$. Since P is an $A \cup \{u\}$-path in Π_1, we have $\alpha \neq 0$. Now note that

$\Pi_1 \in \mathcal{P}(G,A)$, $\mathrm{val}_A(\Pi_1) = \nu(\mathcal{P},A) - 1$, and $p, (u,\alpha) \in B_A(\Pi_1)$. Let $\Pi_2 \in \mathcal{P}^*$ with $(v,0) \in B_A(\Pi_2)$. Applying Lemma 3.3 to Π_1 and Π_2 we find $\Pi' \in \mathcal{P}(G,A)$ with $\mathrm{val}_A(\Pi') = \nu(\mathcal{P},A)$ and either $p \in B_A(\Pi')$ or $(u,\alpha) \in B_A(\Pi')$, as required by the lemma.

Case 2: *There exists* $\Pi_1 \in \mathcal{P}(G, A \cup \{u\})$ *with* $\mathrm{val}_{A\cup\{u\}}(\Pi) = \nu(\mathcal{P},A) + 1$.

There is a path $P \in \Pi_1$ whose start or end is u. If P is a loose path, then $\Pi' := \Pi_1 - \{P\}$ satisfies the lemma. Therefore we may assume that P is an $A \cup \{u\}$-path; furthermore, by possibly reversing P, we may assume that u is the end of P. Let $\alpha = \omega_G(P)$. Since P is an $A \cup \{u\}$-path in Π_1, we have $\alpha \neq 0$. Now note that Π_1 is an A-collection, $\mathrm{val}_A(\Pi_1) = \nu(\mathcal{P},A)$, and $(u,\alpha) \in B_A(\Pi_1)$. Thus $\Pi' := \Pi_1$ satisfies the lemma.

Case 3: *There exists* $\Pi_1 \in \mathcal{P}(G, A \cup \{u\})$ *with* $\mathrm{val}_{A\cup\{u\}}(\Pi_1) = \nu(\mathcal{P},A)$ *and there exists* $(z,\beta) \in B_{A\cup\{u\}}(\Pi_1) - \{(w,\delta)\}$ *with* $zu = f \in E(G)$.

Let $P \in \Pi_1$ be the path ending at w. We may assume that u is the start of P, since otherwise we reduce to Case 1. Let $P_z \in \Pi_1$ be the path ending at z, let $P_u = (P_z, f, u)$, and let $P_w = (P_z, f, P)$. Let $\alpha = \omega_G(P_u)$. Note that $\omega_G(P_w) = \alpha + \delta$, so either $\alpha \neq 0$ or $\omega_G(P_w) = \delta$. Suppose that $\alpha \neq 0$. Let $\Pi' = (\Pi_1 - \{P, P_z\}) \cup \{P_u\}$. Note that $\mathrm{val}_A(\Pi') = \nu(\mathcal{P},A)$ and $(u,\alpha) \in B_A(\Pi')$, as required. Now suppose that $\omega_G(P_w) = \delta$. Let $\Pi' = (\Pi_1 - \{P, P_z\}) \cup \{P_w\}$. Note that $\mathrm{val}_A(\Pi') = \nu(\mathcal{P},A)$ and $(w,\delta) \in B_A(\Pi')$, as required.

Case 4: *There exists* $\Pi_2 \in \mathcal{P}(G, A \cup \{u\})$ *such that* $\mathrm{val}_{A\cup\{u\}}(\Pi_2) = \nu(\mathcal{P},A)$ *and* $(z,\beta), (v,0) \in B_{A\cup\{u\}}(\Pi_2)$ *where* $zu = f \in E(G)$ *and* $(z,\beta) \notin \{(w,\delta),(v,0)\}$.

Note that, since $(v,0) \in \mathcal{R}(\mathcal{P},A)$, we have $(v,0) \neq (w,\delta)$. Recall that $\Pi \in \mathcal{P}(G, A \cup \{u\})$, $\mathrm{val}_{A\cup\{u\}}(\Pi) = \nu(\mathcal{P},A)$, and $(w,\delta) \in B_{A\cup\{u\}}(\Pi)$. Applying Lemma 3.4 to $\Pi_1 := \Pi$ and Π_2, we find $\Pi_3 \in \mathcal{P}(G,A)$ such that either $\mathrm{val}_{A\cup\{u\}}(\Pi_3) > \nu(\mathcal{P},A)$, or $\mathrm{val}_{A\cup\{u\}}(\Pi_3) = \nu(\mathcal{P},A)$ and either $(v,0),(w,\delta) \in B_{A\cup\{u\}}(\Pi_3)$ or $(z,\beta),(w,\delta) \in B_{A\cup\{u\}}(\Pi_3)$. The case that $\mathrm{val}_{A\cup\{u\}}(\Pi_3) > \nu(\mathcal{P},A)$ reduces to Case 2 and the case that $\mathrm{val}_{A\cup\{u\}}(\Pi_3) = \nu(\mathcal{P},A)$ and either $(v,0),(w,\delta) \in B_{A\cup\{u\}}(\Pi_3)$ or $(z,\beta),(w,\delta) \in B_{A\cup\{u\}}(\Pi_3)$ reduces to Case 3.

Case 5: *There exists* $\Pi_2 \in \mathcal{P}(G, A \cup \{u\})$ *such that* $\mathrm{val}_{A\cup\{u\}}(\Pi_2) = \nu(\mathcal{P},A)$ *and* $(u,0),(v,0) \in B_{A\cup\{u\}}(\Pi_2)\}$.

Note that, since $(v,0) \in \mathcal{R}(\mathcal{P},A)$, we have $(v,0) \neq (w,\delta)$. Recall that $\Pi \in \mathcal{P}(G, A \cup \{u\})$, $\mathrm{val}_{A\cup\{u\}}(\Pi) = \nu(\mathcal{P},A)$, and $(w,\delta) \in B_{A\cup\{u\}}(\Pi)$. Applying Lemma 3.4 to $\Pi_1 := \Pi$ and Π_2, we find $\Pi_3 \in \mathcal{P}(G,A)$ such that either $\mathrm{val}_{A\cup\{u\}}(\Pi_3) > \nu(\mathcal{P},A)$, or $\mathrm{val}_{A\cup\{u\}}(\Pi_3) =$

$\nu(\mathcal{P}, A)$ and either $(u, 0), (w, \delta) \in B_{A\cup\{u\}}(\Pi_3)$ or $(v, 0), (w, \delta) \in B_{A\cup\{u\}}(\Pi_3)$. The case that $\mathrm{val}_{A\cup\{u\}}(\Pi_3) > \nu(\mathcal{P}, A)$ reduces to Case 2; the case that $\mathrm{val}_{A\cup\{u\}}(\Pi_3) = \nu(\mathcal{P}, A)$ and $(v, 0), (w, \delta) \in B_{A\cup\{u\}}(\Pi_3)$ reduces to Case 3; and the case that $\mathrm{val}_{A\cup\{u\}}(\Pi_3) = \nu(\mathcal{P}, A)$ and $(u, 0), (w, \delta) \in B_{A\cup\{u\}}(\Pi_3)$ reduces to Case 1.

Let $\Pi_v \in \mathcal{P}^*$ with $(v, 0) \in B_A(\Pi_v)$. We may assume that there is a path $P \in \Pi_v$ that contains u, since otherwise $\Pi_2 := \Pi_v \cup \{(u)\}$ meets the criteria of Case 5.

Case 6: *P is a loose path with respect to A.*

For any $y \in V(P)$, we let P_y denote the initial segment of P ending at y. We may assume that $\omega_G(P_u) = 0$, since otherwise $\Pi' := (\Pi_v - \{P\}) \cup \{P_u\}$ satisfies the lemma. Now we may assume that v is the end of P, since otherwise $\Pi_2 := ((\Pi_v - \{P\}) \cup \{P_u\}$ meets the criteria of Case 5. Now let z be the vertex preceding u on P and let P' be the subpath of P starting at u and ending at v. Let $\beta = \omega_G(P_z)$. We may assume that $(z, \beta) \neq (w, \delta)$, since otherwise $\Pi' := ((\Pi_v - \{P\}) \cup \{P_z\}$ satisfies the lemma. Finally, we see that $\Pi_2 := (\Pi_v - \{P\}) \cup \{P_z, P'\}$ meets the criteria of Case 4.

Case 7: *P is an A-path.*

For any $y \in V(P)$, we let P_y denote the initial segment of P ending at y. Note that, by possibly reversing the direction of P, we may assume that $\omega_G(P_u) \neq 0$; let $\alpha = \omega_G(P_u)$. Let z be the vertex on P immediately following u, let P' denote the subpath of $\bar{P}$ that ends at z, and let $\beta = \omega_G(P')$. We may assume that $(z, \beta) = (w, \delta)$, since otherwise $\Pi_2 := (\Pi_v - \{P\}) \cup \{P_u, P'\}$ meets the criteria of Case 4. Let $Q \in \Pi_v$ be the path ending at v. Let $P'' = (Q, e, \bar{P}_u)$. Note that P'' is an A-path and that $\omega_G(P'') = \omega_G(Q) + \omega_G(e, u) - \omega_G(P_u) = -\alpha \neq 0$. Therefore $\Pi' := (\Pi_v - \{P, Q\}) \cup \{P', P''\}$ satisfies the lemma.

That completes the proof; this proof can easily be made algorithmic with the desired running time. □

We say that (G, A) is *$\mathcal{P}$-critical* if G is connected, $E_0(G, A) = \emptyset$, $D_1(\mathcal{P}, A) = A$, and $D_2(\mathcal{P}, A) = V(G) - A$. The next result generalizes Lemma 2.7.

Lemma 4.5. *If (G, A) is $\mathcal{P}$-critical and $\mathrm{def}(\mathcal{P}, A) > 1$, then there exists $\Pi \in \mathcal{P}(G, A)$ such that $\mathrm{val}_A(G) = \nu_A(\mathcal{P}) + 1$. Moreover, if $|\mathcal{P}| \leq 2|V(G)|$, then we can find such Π' in $O(|V(G)|^4)$ time.*

Proof. We start by considering an easy case.

Case 1: *There exists $\Pi_1 \in \mathcal{P}(G, A)$ with $\mathrm{val}_A(\Pi_1) = \nu_A(\mathcal{P})$ and there exists $(u, \alpha), (v, \beta) \in B_A(\Pi_1)$ where $uv = e \in E(G)$.*

We break this into two further subcases.

Case 1.1: $\alpha + \omega_G(e, v) - \beta \neq 0$.

Let $P_u, P_v \in \Pi_1$ be the paths ending at u and v respectively and let $P = (P_u, e, \bar{P}_v)$. Note that P is an A-path and that $\omega_G(P) = \alpha + \omega_G(e, v) - \beta \neq 0$. Thus $\Pi := (\Pi_1 - \{P_u, P_v\}) \cup \{P\}$ satisfies the lemma.

Case 1.2: $\alpha + \omega_G(e, v) - \beta = 0$.

Note that, since (G, A) is $\mathcal{P}$-critical, either $u \notin A$ or $v \notin A$. By possibly swapping u and v, we may assume that $v \notin A$. Then, since (G, A) is $\mathcal{P}$-critical, there exists $\beta' \in \Gamma(\mathcal{P}, A, v) - \{\beta\}$. Let $\Pi_2 \in \mathcal{P}^*$ with $(v, \beta') \in B_A(\Pi_2)$. Applying Lemma 3.4 to Π_2 and Π_1, we find $\Pi_3 \in \mathcal{P}(G, A)$ such that either $\mathrm{val}_A(\Pi_3) > \nu_A(\mathcal{P})$ or $\mathrm{val}_A(\Pi_3) = \nu_A(\mathcal{P})$ and either $(u, \alpha), (v, \beta') \in B_A(\Pi_3)$ or $(v, \beta), (v, \beta') \in B_A(\Pi_3)$. If $\mathrm{val}_A(\Pi_3) > \nu_A(\mathcal{P})$, then $\Pi := \Pi_3$ satisfies the lemma. Also, note that $B_A(\Pi_3)$ cannot contain both (v, β) and (v, β'). Therefore we may assume that $(u, \alpha), (v, \beta') \in B_A(\Pi_3)$. Now, since $\beta \neq \beta'$, we have $\alpha + \omega_G(e, v) - \beta' \neq \alpha + \omega_G(e, v) - \beta = 0$. Therefore $\Pi_1 := \Pi_3$ satisfies the criterion for Case 1.1.

$(*)$ Among all triples $(\Pi_1, (v_1, \alpha_1), (v_2, \alpha_2))$ where $\Pi_1 \in \mathcal{P}(G, A)$, $\mathrm{val}_A(\Pi_1) = \nu_A(\mathcal{P})$, and $(v_1, \alpha_1), (v_2, \alpha_2) \in B_A(\Pi_1)$ we choose the triple such that the distance between v_1 and v_2 in G is minimum.

In view of Case 1, we may assume that v_1 is not adjacent to v_2. Let P be a shortest (v_1, v_2)-path and let u be an internal vertex of P. Since (G, A) is $\mathcal{P}$-critical, there exists $\beta \in \Gamma(\mathcal{P}, A, u)$. Let $\Pi_2 \in \mathcal{P}^*$ with $(u, \beta) \in B_A(\Pi_2)$. Applying Lemma 3.4 to Π_2 and Π_1, we find $\Pi_3 \in \mathcal{P}(G, A)$ such that either $\mathrm{val}_A(\Pi_3) > \nu_A(\mathcal{P})$ or $\mathrm{val}_A(\Pi_3) = \nu_A(\mathcal{P})$ and $(u, \beta), (v_i, \alpha_i) \in B_A(\Pi_3)$ for some $i \in \{1, 2\}$. If $\mathrm{val}_A(\Pi_3) > \nu_A(\mathcal{P})$, then $\Pi := \Pi_3$ satisfies the lemma. Thus, by symmetry, we may assume that $\mathrm{val}_A(\Pi_3) = \nu_A(\mathcal{P})$ and $(u, \beta), (v_1, \alpha_1) \in B_A(\Pi_3)$. However, since v_1 is closer to u than it is to v_2, we have a contradiction to $(*)$.

That completes the proof; this proof can easily be made algorithmic with the desired running time. □

5. The algorithm

Throughout the algorithm we maintain a set $\mathcal{P} \subseteq \mathcal{P}(G, A)$. We are primarily interested in the sets $D_1(\mathcal{P}, A)$ and $D_2(\mathcal{P}, A)$. Therefore, by removing unnecessary A-collections from $\mathcal{P}$, we keep

$$|\mathcal{P}| \leq |D_1(\mathcal{P}, A)| + 2|D_2(\mathcal{P}, A)| \leq 2|V(G)|.$$

If $\mathcal{P}_1, \mathcal{P}_2 \subseteq \mathcal{P}(G, A)$, then we say that $\mathcal{P}_2$ is *richer* than $\mathcal{P}_1$, with respect to A, if either $\nu_A(\mathcal{P}_2) > \nu_A(\mathcal{P}_1)$ or $\nu_A(\mathcal{P}_2) = \nu_A(\mathcal{P}_1)$ and $|D_1(\mathcal{P}_2, A)| + 2|D_2(\mathcal{P}_2, A)| > |D_1(\mathcal{P}_1, A)| + 2|D_2(\mathcal{P}_1, A)|$.

By possibly shifting (as we did in Lemma 1.2), we may assume that (G, A) satisfies:

(1) for each $v \in D_1(\mathcal{P}, A)$, $\Gamma(\mathcal{P}, A, v) = \{0\}$, and
(2) for each $u \in N_G(D(\mathcal{P}, A)) - A$, there exists $uv = e \in E(G)$ such that $\omega_G(e, v) \in \Gamma(\mathcal{P}, A, v)$.

Now let $A' = A \cup N_G(D(\mathcal{P}, A)) \cup D_1(\mathcal{P}, A)$ and $X = N_{G-E_0(G,A')}(D(\mathcal{P}, A))$.

Optimality condition: *If* $\mathrm{def}(\mathcal{P}, A) = \mathrm{odd}(G-X, A'-X) - |A'-A| - |X|$, *then the A-collections in $\mathcal{P}^*$ are optimal.*

Proof. Note that $\mathrm{def}(\mathcal{P}, A) \geq \mathrm{def}(G, A) \geq \mathrm{odd}(G - X, A' - X) - |A' - A| - |X|$. Thus, if $\mathrm{def}(\mathcal{P}, A) = \mathrm{odd}(G - X, A' - X) - |A' - A| - |X|$, then $\mathrm{def}(\mathcal{P}, A) = \mathrm{def}(G, A)$ and, hence, each A-collection in $\mathcal{P}^*$ is optimal. □

In each iteration of the algorithm, if $\mathrm{def}(\mathcal{P}, A) \neq \mathrm{odd}(G - X, A' - X) - |A' - A| - |X|$, then we find an A-collection Π such that $\mathcal{P} \cup \{\Pi\}$ is richer than $\mathcal{P}$. Hence in at most $O(|V(G)|^2)$ iterations we will find an optimal A-collection. It remains to show how we find the promised A-collection Π.

We omit the elementary proof of the next lemma.

Lemma 5.1. *Let $A_1, X_1 \subseteq V(G)$ such that $A \cup X_1 \subseteq A_1 \subseteq A'$ and $X_1 \subseteq X$. Then, in $O(|V(G)|^3)$ time, we can construct $\mathcal{P}_1 \subset \mathcal{P}(G - X_1, A_1 - X_1)$ such that either $\nu_{A_1}(\mathcal{P}_1) > \nu_A(\mathcal{P}) - |X_1|$ or $\nu_{A_1}(\mathcal{P}_1) = \nu_A(\mathcal{P}) - |X_1|$ and $\mathcal{R}(\mathcal{P}, A) \subseteq \mathcal{R}(\mathcal{P}_1, A_1)$.*

Lemma 5.2. *Let $A'' \subseteq A'$ with $A \subseteq A''$. Suppose that $\Pi' \in \mathcal{P}(G, A'')$ where either*

(i) *$\nu_{A''}(\mathcal{P}') > \nu_A(\mathcal{P})$ or*
(ii) *$\nu_{A''}(\mathcal{P}') = \nu_A(\mathcal{P})$ and there exists $(v, \beta) \in B_{A''}(\Pi') - \mathcal{R}(\mathcal{P}, A)$ where $v \notin D_2(\mathcal{P}, A)$.*

Then, in $O(|V(G)|^4)$ time, we can find $\Pi \in \mathcal{P}(G, A)$ such that $\mathcal{P} \cup \{\Pi\}$ is richer than $\mathcal{P}$.

Proof. The proof is inductive on $|A'' - A|$. If $A'' = A$, then $\Pi := \Pi'$ satisfies the lemma. Thus we may assume that there exists $a \in A'' - A$. Let $A_1 = A'' - \{a\}$. By Lemma 5.1, we can construct $\mathcal{P}_1 \subset \mathcal{P}(G, A_1)$ such that either $\nu_{A_1}(\mathcal{P}_1) > \nu_A(\mathcal{P})$ or $\nu_{A_1}(\mathcal{P}_1) = \nu_A(\mathcal{P})$ and $\mathcal{R}(\mathcal{P}, A) \subseteq \mathcal{R}(\mathcal{P}_1, A_1)$. Inductively, we may assume that $\nu_{A_1}(\mathcal{P}_1) =$

$\nu_A(\mathcal{P})$, $D_1(\mathcal{P}_1, A_1) = D_1(\mathcal{P}, A)$, and $D_2(\mathcal{P}_1, A_1) = D_2(\mathcal{P}, A)$. Now, by Lemma 4.4, we can construct $\Pi'' \in \mathcal{P}(G, A_1)$ such that $\mathcal{P}_1 \cup \{\Pi''\}$ is richer than $\mathcal{P}_1$ with respect to A_1. □

The next lemma is proved similarly; we leave the details to the reader.

Lemma 5.3. *Let $X' \subseteq X$. Suppose that $\Pi' \in \mathcal{P}(G - X', A' - X')$ where either*

(i) $\nu_{A'-X'}(\mathcal{P}') > \nu_A(\mathcal{P})$ *or*
(ii) $\nu_{A'-X'}(\mathcal{P}') = \nu_A(\mathcal{P})$ *and there exists* $(v, \beta) \in B_{A'-X'}(\Pi') - \mathcal{R}(\mathcal{P}, A)$ *where* $v \notin D_2(\mathcal{P}, A)$.

Then, in $O(|V(G)|^4)$ time, we can find $\Pi \in \mathcal{P}(G, A)$ such that $\mathcal{P} \cup \{\Pi\}$ is richer than $\mathcal{P}$.

Let $G_1 = G - X$ and let $A_1 = A' - X$. Now, by Lemma 5.1, we can construct $\mathcal{P}_1 \subset \mathcal{P}(G_1, A_1)$ such that either $\nu_{A_1}(\mathcal{P}_1) > \nu_A(\mathcal{P})$ or $\nu_{A_1}(\mathcal{P}_1) = \nu_A(\mathcal{P})$ and $\mathcal{R}(\mathcal{P}, A) \subseteq \mathcal{R}(\mathcal{P}_1, A_1)$. By Lemma 5.3, we may assume that $\nu_{A_1}(\mathcal{P}_1) = \nu_A(\mathcal{P})$, $D_1(\mathcal{P}_1, A_1) = D_1(\mathcal{P}, A)$, and $D_2(\mathcal{P}_1, A_1) = D_2(\mathcal{P}, A)$. Now let $G_2 = G_1 - E_0(G_1, A_1)$. Note that no A_1-collection in G_1 uses an edge in $E_0(G_1, A_1)$, so $P_1 \subseteq \mathcal{P}(G_2, A_1)$. Note that, if we can find $\Pi' \in \mathcal{P}(G_2, A_1)$ such that $\mathrm{val}_{A_1}(\Pi') > \nu_{A_1}(\mathcal{P}_1)$, then, by Lemma 5.3, we can construct $\Pi \in \mathcal{P}(G, A)$ such that $\mathcal{P} \cup \{\Pi\}$ is richer than $\mathcal{P}$.

Let H be a component of G_2, let $A_H = A_1 \cap V(H)$. For each $\Pi \in \mathcal{P}(G_2, A_1)$, we let $\Pi|H$ denote the restriction of Π to H and let $\Pi - H$ denote the restriction of Π to $G_2 - H$. Let $\Pi_1, \Pi_2 \in \mathcal{P}_1^*$. Suppose that $\mathrm{val}_{A_H}(\Pi_1|H) > \mathrm{val}_{A_H}(\Pi_2|H)$. Now let $\Pi' = (\Pi_2 - H) \cup (\Pi_1|H)$. Note that $\Pi' \in \mathcal{P}(G_2, A_1)$ and that $\mathrm{val}_{A_1}(\Pi') > \nu_{A_1}(\mathcal{P}_1)$, as required. Therefore we may assume that, for all $\Pi_1, \Pi_2 \in \mathcal{P}_1^*$, we have $\mathrm{val}_{A_H}(\Pi_1|H) > \mathrm{val}_{A_H}(\Pi_2|H)$. Let $\mathcal{P}_H = \{\Pi|H \,:\, \Pi \in \mathcal{P}_1^*\}$.

Lemma 5.4. *For each component H of G_2, either* $\mathrm{def}(H, A_H) = 0$ *or* (H, A_H) *is $\mathcal{P}_H$-critical.*

Proof. Note that, if $uv = e \in E(G)$ with $u \in V(G) - D(\mathcal{P}, A)$ and $v \in D(\mathcal{P}, A)$, then either $u \in X$ or $e \in E_0(G, A')$. Moreover, $D(\mathcal{P}_1, A_1) = D(\mathcal{P}, A)$. Thus, if H is a component of G_2, then either $V(H) \subseteq D(\mathcal{P}_1, A_1)$ or $V(H) \cap D(\mathcal{P}_1, A_1) = \emptyset$. If $V(H) \cap D(\mathcal{P}_1, A_1) = \emptyset$, then $\mathrm{def}(\mathcal{P}_H, A_H) = 0$. Thus we may assume that $V(H) \subseteq D(\mathcal{P}_1, A_1)$. Note that, since H is a component of G_2, $D_1(\mathcal{P}_H, A_H) = D_1(\mathcal{P}_1, A_1) \cap V(H)$ and $D_2(\mathcal{P}_H, A_H) = D_2(\mathcal{P}_1, A_1) \cap V(H)$. By the definition of A', a vertex $v \in D(\mathcal{P}_1, A_1)$ is in $D_1(\mathcal{P}_1, A_1)$ if and only if $v \in A_1$. Hence H is $\mathcal{P}_H$-critical, as required. □

Suppose that (H, A_H) is $\mathcal{P}_H$-critical and that $\text{def}(\mathcal{P}_H, A_H) > 1$. Then, by Lemma 4.5, we can construct $\Pi_1 \in \mathcal{P}(H, A_H)$ such that $\text{val}_{A_H}(\Pi_1) > \nu(\mathcal{P}_H, A_H)$. Now let $\Pi_2 \in \mathcal{P}_1^*$ and let $\Pi' = \Pi_1 \cup (\Pi_2 - H)$. Note that $\Pi' \in \mathcal{P}(G_2, A_1)$ and that $\text{val}_{A_1}(\Pi') > \nu_{A_1}(\mathcal{P}_1)$, as required. Therefore we may assume that: *For each component H of G_2, we have* $\text{def}(\mathcal{P}_H, A_H) \leq 1$. Thus $\text{def}(G_1, A_1) = \text{odd}(G_1, A_1)$. So, we have:

$$\begin{aligned} \text{def}(\mathcal{P}, A) &= \text{def}(G - X, A' - X) - |A' - A| - |X| \\ &= \text{odd}(G - X, A' - X) - |A' - A| - |X|, \end{aligned}$$

as required. This completes the description and proof of the algorithm.

Let $n = |V(G)|$. The algorithm, as stated, requires $O(n^6)$ time. The complexity in Lemma 3.2 can be improved from $O(n^3)$ to $O(n^2)$, by combining the proofs of Lemma 3.2 and 3.1. This would reduce the overall complexity of our algorithm from $O(n^6)$ to $O(n^5)$.

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